

NAME

lock, unlock, query_lock, set_lock_cache_enable, lock_cache_enabled – Class ss_m Methods for Locking

SYNOPSIS

```
#include <sm_vas.h> // which includes sm.h

static rc_t lock(
    const lvid_t&      lvid,
    lock_mode_t        mode,
    lock_duration_t    duration = t_long,
    long               timeout = WAIT_SPECIFIED_BY_XCT);

static rc_t lock(
    const lockid_t&    lockid,
    lock_mode_t        mode,
    lock_duration_t    duration = t_long,
    long               timeout = WAIT_SPECIFIED_BY_XCT);

static rc_t unlock(const lockid_t& lockid);

static rc_t query_lock(
    const lockid_t&    lockid,
    lock_mode_t        mode,
    bool               implicit = false);

static rc_t set_lock_cache_enable(bool enable);
static rc_t lock_cache_enabled(bool& enabled);

static rc_t set_escalation_thresholds(
    int4             toPage,
    int4             toStore,
    int4             toVolume);

static rc_t get_escalation_thresholds(
    int4&            toPage,
    int4&            toStore,
    int4&            toVolume);

static rc_t dont_escalate(
    const lockid_t&    n,
    bool               passOnToDescendants = true);
static rc_t dont_escalate(
    const lvid_t&      lvid,
    bool               passOnToDescendants = true);

// Lock ID

class lockid_t {
public:
    //
    // The lock graph consists of 6 node: volumes, stores, pages, key values,
    // records, and extents. The first 5 of these form a tree of 4 levels.
```

```

// The node for extents is not connected to the rest.
// The node_space_t enumerator maps node types to integers.
// These numbers are used for
// indexing into arrays containing node type specific info per entry (e.g
// the lock caches for volumes, stores, and pages).
//
enum name_space_t {
t_bad      = 10,
t_vol      = 0,
t_store    = 1,    // parent is 1/2 = 0 t_vol
t_page     = 2,    // parent is 2/2 = 1 t_store
t_kv1      = 3,    // parent is 3/2 = 1 t_store
t_record   = 4,    // parent is 4/2 = 2 t_page
t_extent   = 5,
t_user1    = 6,
t_user2    = 7,    // parent is t_user1
t_user3    = 8,    // parent is t_user2
t_user4    = 9     // parent is t_user3
};

struct user1_t {
uint2_t      u1;
    user1_t() : u1(0) {}
    user1_t(uint2_t v1) : u1(v1) {}
};

struct user2_t : public user1_t {
uint4_t      u2;
    user2_t() : u2(0) {}
    user2_t(uint2_t v1, uint4_t v2): user1_t(v1), u2(v2) {}
};

struct user3_t : public user2_t {
uint4_t      u3;
    user3_t() : u3(0) {}
    user3_t(uint2_t v1, uint4_t v2, uint4_t v3)
        : user2_t(v1, v2), u3(v3) {}
};

struct user4_t : public user3_t {
uint4_t      u4;
    user4_t() : u4(0) {}
    user4_t(uint2_t v1, uint4_t v2, uint4_t v3, uint4_t v4)
        : user3_t(v1, v2, v3), u4(v4) {}
};

bool operator==(const lockid_t& p) const;
bool operator!=(const lockid_t& p) const;
friend ostream& operator<<(ostream& o, const lockid_t& i);

uint4_t      hash() const;
void         zero();

name_space_t lspace() const;

```

```

vid_t          vid() const;
const snum_t&   store() const;
const extnum_t& extent() const;
const shpid_t&  page() const;
const slotid_t& slot() const;
uint2_t        u1() const;
uint4_t        u2() const;
uint4_t        u3() const;
uint4_t        u4() const;

void           set_ext_has_page_alloc(bool value);
bool           ext_has_page_alloc() const ;

NORET          lockid_t() ;
NORET          lockid_t(const vid_t& vid);
NORET          lockid_t(const extid_t& extid);
NORET          lockid_t(const stid_t& stid);
NORET          lockid_t(const lpid_t& lpid);
NORET          lockid_t(const stpgid_t& stpgid);
NORET          lockid_t(const rid_t& rid);
NORET          lockid_t(const kvl_t& kvl);
NORET          lockid_t(const lockid_t& i);

NORET          lockid_t(const user1_t& u);
NORET          lockid_t(const user2_t& u);
NORET          lockid_t(const user3_t& u);
NORET          lockid_t(const user4_t& u);

void           extract_extent(extid_t &e) const;
void           extract_stid(stid_t &s) const;
void           extract_lpid(lpid_t &p) const;
void           extract_rid(rid_t &r) const;
void           extract_kvl(kvl_t &k) const;
void           extract_user1(user1_t &u) const;
void           extract_user2(user2_t &u) const;
void           extract_user3(user3_t &u) const;
void           extract_user4(user4_t &u) const;

bool           IsUserLock() const;

void           truncate(name_space_t space);

lockid_t&      operator=(const lockid_t& i);

};

ostream& operator<<(ostream& o, const lockid_t::user1_t& u);
ostream& operator<<(ostream& o, const lockid_t::user2_t& u);
ostream& operator<<(ostream& o, const lockid_t::user3_t& u);
ostream& operator<<(ostream& o, const lockid_t::user4_t& u);

istream& operator>>(istream& o, lockid_t::user1_t& u);
istream& operator>>(istream& o, lockid_t::user2_t& u);
istream& operator>>(istream& o, lockid_t::user3_t& u);

```

```
istream& operator>>(istream& o, lockid_t::user4_t& u);
```

DESCRIPTION

Locks are acquired implicitly by many **ss_m** methods. For those situations where more precise control of locking is desired, the following methods allow explicit locking and unlocking.

The class representing a generic lock is a `lockid_t`, described above. The SSM acquires locks on pages, extents, records, stores, and volumes. **The extent locks are NOT to be used by VASs**, simply because the extent-based structure of the SSM is likely to change in future releases.

lock(lvid, mode, duration, timeout)

lock(lockid, mode, duration, timeout)

The **lock** method is used to acquire a lock on volume, index, file or record. The first version of the method locks the volume specified by *lvid*. The second version locks the index, file or record specified by *lockid*. The *mode* parameter specifies the lock mode to acquire. Valid `lock_mode_t` values are listed in

basics.h. The *duration* parameter specifies how long the lock will be held. Valid values (among those listed in **basics.h**) are: `t_instant`, `t_short` and `t_long`. The *timeout* parameter specifies how long to wait for a lock.

unlock(lockid)

The **unlock** method releases the most recently acquired lock on the file, index, or record identified by *lockid*. Note, that only locks with duration **t_short** can be released before end-of-transaction.

query_lock(lockid, mode, implicit)

The **query_lock** method the mode of the lock held on *lockid* by the current transaction. The lock mode is returned in *mode* and will be **NL** (no lock) if not locked. If *implicit* is **false** then only explicit locks on *lockid* will be considered. For example, if file F is **SH** locked and a query is made about a record in F, the mode returned will be **NL**. **However, if *implicit* is **true**, then **SH** would be returned for this example.**

Lock Cache Control

Each transaction has a cache of recently acquired locks. The following methods control the use of the cache. These are not supported methods and may be removed in later versions of the software. Note: that the methods only affect the transaction associated with the current thread.

set_lock_cache_enable(enable)

The **set_lock_cache_enable** method turns on the cache if *enable* is **true** and turns it off otherwise.

lock_cache_enabled(enabled)

The **lock_cache_enabled** method sets *enabled* to **true** if the lock cache is on.

Escalation

The lock manager will escalate from a record lock to a page lock, from a page lock to a store lock, and from a store lock to a volume lock, to reduce the number of locks in the table. You can control the thresholds for escalation through the methods **get_escalation_thresholds** and **set_escalation_thresholds**. The default values are as follows:

record-to-page
5
page-to-store
25
store-to-volume
0

In all cases, a threshold of 0 prevents escalation.

When escalation is in use, it be prevented on selected volumes or other lock-able objects through the three **dont_escalate** methods. If the argument *passOnToDescendants* is *false*, locks acquired on objects below the volume (or given lockid) in the lock hierarchy will still be escalated according to the thresholds.

ERRORS

TODO

EXAMPLES

TODO

VERSION

This manual page applies to Version 2.0 of the Shore Storage Manager.

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SEE ALSO

transaction(ssm), **id(ssm)**, and **intro(ssm)**.