## YOU SHOULD ATTEMPT ALL 4 PROBLEMS

1. A positive number $n$ is called squarefree if no square divides it except for 1. Give an $O(n)$ algorithm that lists all the squarefree numbers in $\{1, \ldots, n\}$. You can treat arithmetic operations such as addition or subtraction or multiplication or integer division as operations with unit cost.
2. Suppose there exists a randomized polynomial-time algorithm with the following behavior: On input an unsatisfiable Boolean formula, it outputs "no" with probability at least $3 / 4$; on input a Boolean formula with exactly one satisfying assignment, it outputs "yes" with probability at least $3 / 4$; on any other input, the behavior of the algorithm is unspecified. Prove that the existence of such an algorithm implies that the polynomial-time hierarchy collapses. To which level can you prove the PH collapses to?
3. A pair of vertices $s$ and $t$ in a graph $G$ is called 2-connected if there are at least two edge-disjoint paths between $s$ and $t$ in the graph. Given a connected graph $G=(V, E)$ with weights $w_{e}$ on edges and vertices $s$ and $t$, a 2-route $s$-t cut is a set of edges $E^{\prime} \subset E$ such that upon removing $E^{\prime}$ from the graph, $s$ and $t$ are not 2-connected. The weight of a cut $E^{\prime}$ is given by $w\left(E^{\prime}\right)=\sum_{e \in E^{\prime}} w_{e}$.
a. Prove that $s$ and $t$ are 2 -connected if and only if every $s$ - $t$ cut in the graph contains at least 2 edges.
b. Use part (a) to give a polynomial time algorithm for finding the minimum weight 2 -route $s-t$ cut in a given graph.
4. In honor of the upcoming election, consider the problem of apportioning seats in a legislature. Suppose there are $m$ states, with populations $p_{1}, \ldots, p_{m}$, and the legislature has $n$ seats. (Currently for the U.S. Congress, $m=50$ and $n=435$.) If the numbers

$$
q_{i}:=n \frac{p_{i}}{\sum_{i} p_{i}}
$$

were all integers, we would give the $i$-th state $q_{i}$ seats. However, they usually are not, so some adjustment is necessary.

Alexander Hamilton suggested the following rule: round each $q_{i}$ down to the closest integer, and provisionally give the $i$-th state a number of seats equal to this integer. The remaining seats are to be assigned, one each, to the states with the largest values of $q_{i}$.
a. Explain why this problem can be solved in $O(m)$ steps. For this problem, a "step" means an arithmetic operation (addition, subtraction, multiplication, division with remainder), or a comparison, on integers.
b. What problems (such as unfairness from certain perspective, or other undesirable features) might arise by the use of Hamilton's rule?

## Please be concise.

