Microprocessors: 20 Years Back, 10 Years Ahead

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Outline

- The enabler: semiconductor technology
- Role of the processor architect
- Micro-architectures of the past 20 years
 - From pipelining to speculation
- Micro-architectures of the next 10 years



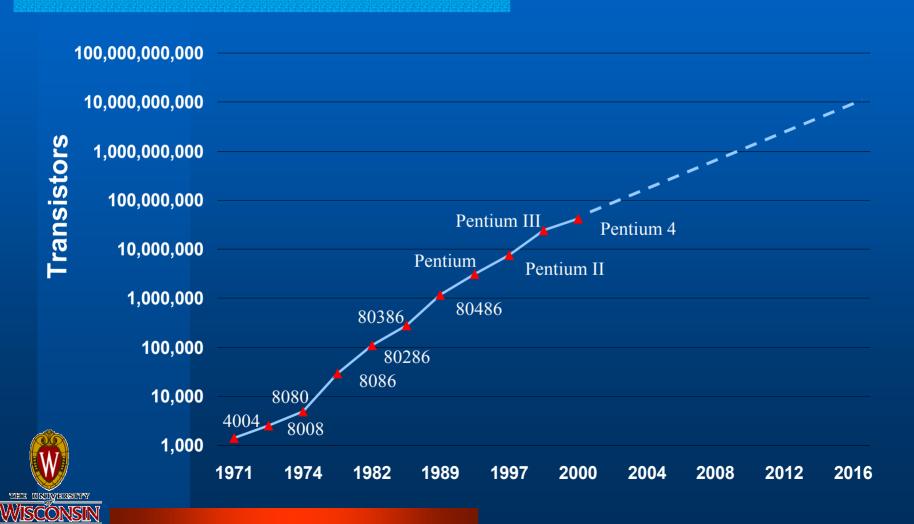
Semiconductor Technology

- Many more available transistors
- Imbalances due to disparate rates of performance improvement
 - E.g., logic and memory speeds

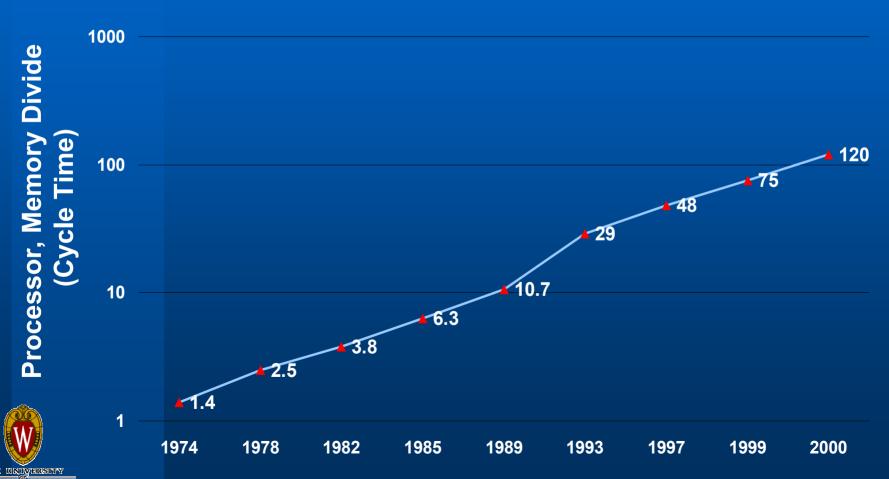
How does this impact the architecture of microprocessors?



Number of Transistors

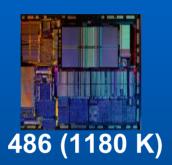


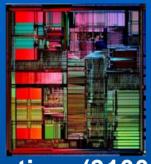
Relative Memory Speed



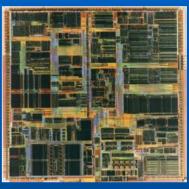
Intel Microprocessors



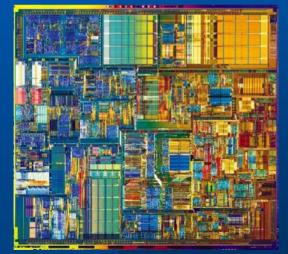




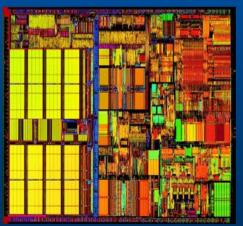
Pentium (3100 K)



Pentium II (7500 K)



Pentium 4 (42000 K)



Pentium III (24000 K)

What is being done with all the transistors?



Role of Computer Architect

- Get desired level of performance
- Determine functionality needed
- Determine how functionality should be implemented



Role of Computer Architect...

- Defining functionality
 - Functionality to deal with increasing latencies (e.g., caches, wires)
 - Functionality to increase parallelism and its exploitation
- Implementing functionality
 - Balancing various technology parameters
 - Ease of design / verification / testing



The Performance Equation

Time = Number of Instructions x Cycles per Instruction x Clock Cycle Time

- Not much can be done about first term in hardware
 - But, ...
- Logic speed increase decreases 3rd term
 - Watch out for possible increase in 2nd term
- Use micro-architectural innovations to decrease 2nd and 3rd terms
 - Reduce latencies
 - Exploit parallelism

Microarchitectural Functionality

- Functionality to cope with increasing memory latencies
- Functionality to exploit parallelism

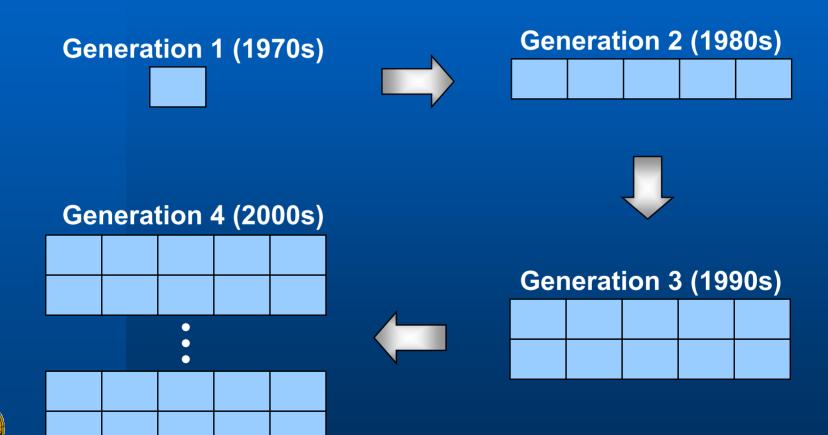


Memory Hierarchies

- Reducing access latency and improving access bandwidth
- Single-level caches
- Multi-level caches
- Non-blocking caches
- Multi-ported and multi-banked caches
- Trace caches



The March of Parallelism



Exploiting Parallelism

- Little change in programming model
 - still write programs in sequential languages
- Automatic parallelization not widely successful
- Great investment in existing software

Resort to low-level, Instruction Level Parallelism (ILP)



Instruction Level Parallelism (ILP)

- Determine small number (e.g., < 100) instructions to be executed
- Determine dependence relationships and create dependence graph
 - Use to determine parallel execution
- Can be done statically (VLIW / EPIC) or dynamically (out-of-order superscalar)



Limitations to ILP

- Branch instructions inhibit determination of instructions to execute: control dependences
- Imperfect analysis of memory addresses inhibits reordering of memory operations: ambiguous memory dependences
- Program/algorithm data flow inhibits parallelism: true dependences
- Increasing latencies exacerbate impact of dependences



Use speculation to overcome impact of dependences

Speculation

Speculation: ".. to assume a business risk in hope of gain"

Webster



Speculation and Computer Architecture

- Speculate outcome of event rather than waiting for outcome to be known
 - Program behavior provides rationale for high success rate
- Functionality to support speculation
- Functionality to speculate better
- Functionality to minimize mis-speculation penalty

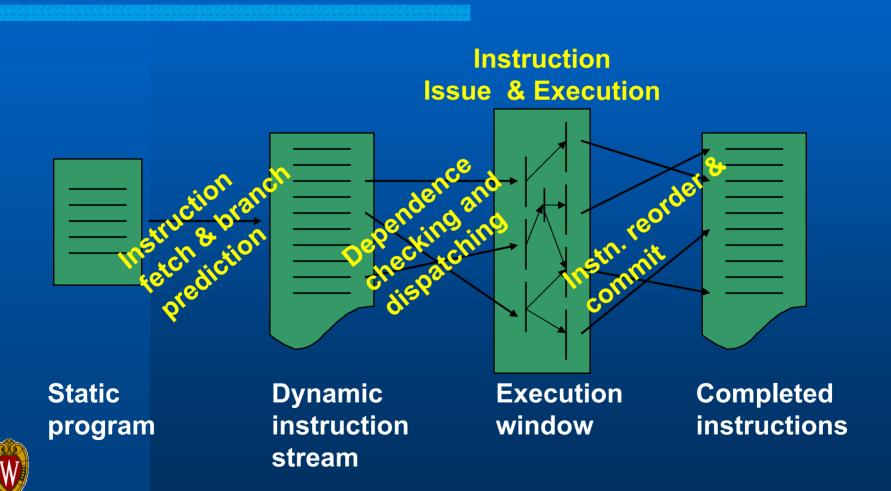


Control Speculation

- Predict outcome of branch instructions
- Speculatively fetch and execute instructions from predicted path
 - Increase available parallelism
- Recover if prediction is incorrect



Model for Speculative Execution



Supporting Control Speculation

- Techniques to predict branch outcome: branch predictors
 - Initiating speculation
 - Improving accuracy of speculation
- Techniques to support speculative execution: reservation stations, register renaming etc.
 - Supporting speculative execution
- Techniques to give appearance of sequential execution: reorder buffers, etc.
 - Doing it transparently



Key observation

Basic mechanisms to support control speculation can support other forms of speculation as well



Performance-Inhibiting Constraints

- Control dependences: inhibit creation of instruction window
 - Use control speculation
- Ambiguous data dependences: inhibit parallelism recognition
 - Use data dependence speculation
- True data dependences: inhibit parallelism
 - Use value speculation
- Common mechanisms may support different forms of speculation
- Different techniques to improve accuracy of speculation

Speculation in Use Today

- Address calculation and translation (especially if 2-step process)
- Cache hit
- Memory ordering violation in multiprocessors
- Load/store dependences



Microprocessors – the next 10 years

- Factor of 30 increase in semiconductor resources
 - How to use it?
- New constraints
 - Power consumption
 - Wire delays
 - Design / verification complexity
- New applications?
 - Throughput-oriented workloads
 - Coarse-grain multithreaded applications



Technology Trends

- Design and verification of large number of transistors becoming unwieldy
- Wires getting relatively slower
 - Short wires for fast clock
 - Implies increase latencies; exploit locality of communication
- Power issues becoming very important



Architect's Role Revisited

- Defining functionality
 - New models needed to further increase parallelism exploitation
- Implementing functionality
 - Becoming a dominating factor?
- Speculation is likely to be the key to overcoming constraints



Implications of Trends

- Implementation considerations will imply computing chips with multiple (replicated?) processing cores
 - "multiprocessor" or "multiprocessor-like" or "multithreaded"
 - Will start out as "logical" replication (e.g., SMT)
 - Will move towards "physical" replication (e.g., CMP)
- How to assign work to multiple processing cores?
 - Independent programs (or threads)
 - Parts of a single program



Throughput-Oriented Processing

- Executing multiple, independent programs on underlying parallel micro-architecture
 - Similar to traditional throughput-oriented multiprocessor
 - Significant engineering challenges, but little in ways of architectural / micro-architectural innovation

Can we use underlying "multiprocessor" to speed up execution of single program?



Parallel Processing of Single Program

- Will the promise of explicit / automatic parallelism come true?
- Will new (parallel) programming languages take over the world?

Don't count on it!



Speculative Parallelization

- Sequential languages aren't going away
- Use speculation to overcome inhibitors to "automatic" parallelization
 - Ambiguous dependences
- Divide program into "speculatively parallel" portions or "speculative threads"



Speculative Threads

- Subject of extensive research today
 - Different speculative parallelization models being investigated



Generic circa 2010 Microprocessor

- 4 8 general-purpose processing engines on chip
 - Used to execute independent programs
 - Explicitly parallel programs (when possible)
 - Speculatively parallel threads
 - Helper threads
- Special-purpose processing units (e.g., DSP functionality)
- Elaborate memory hierarchy
- Elaborate inter-chip communication facilities
- Extensive use of different forms of speculation



Summary

- Semiconductor technology has, and will continue to, give computer architects new opportunities
- Architects have used speculation techniques to overcome performance barriers; will likely continue to do so
- Future microprocessors are going to have capability to execute multiple threads of code
- New models of speculation (e.g., thread-level speculation) will be needed to extract more parallelism

