

**NAME**

lock, unlock, query\_lock, set\_lock\_cache\_enable, lock\_cache\_enabled – Class ss\_m Methods for Locking

**SYNOPSIS**

```
#include <sm_vas.h> // which includes sm.h

static rc_t lock(
    const lvid_t&      lvid,
    lock_mode_t        mode,
    lock_duration_t    duration = t_long,
    long               timeout = WAIT_SPECIFIED_BY_XCT);

static rc_t lock(
    const lockid_t&    lockid,
    lock_mode_t        mode,
    lock_duration_t    duration = t_long,
    long               timeout = WAIT_SPECIFIED_BY_XCT);

static rc_t unlock(const lockid_t& lockid);

static rc_t query_lock(
    const lockid_t&    lockid,
    lock_mode_t        mode,
    bool               implicit = false);

static rc_t set_lock_cache_enable(bool enable);
static rc_t lock_cache_enabled(bool& enabled);

static rc_t set_escalation_thresholds(
    int4 toPage,
    int4 toStore,
    int4 toVolume);

static rc_t get_escalation_thresholds(
    int4& toPage,
    int4& toStore,
    int4& toVolume);

static rc_t dont_escalate(
    const lockid_t&    n,
    bool               passOnToDescendants = true);
static rc_t dont_escalate(
    const lvid_t&      lvid,
    bool               passOnToDescendants = true);

// Lock ID

class lockid_t {
public:
    //
    // The lock graph consists of 6 node: volumes, stores, pages, key values,
    // records, and extents. The first 5 of these form a tree of 4 levels.
```

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// The node for extents is not connected to the rest.
// The node_space_t enumerator maps node types to integers.
// These numbers are used for
// indexing into arrays containing node type specific info per entry (e.g
// the lock caches for volumes, stores, and pages).
//
enum name_space_t {
t_bad      = 10,
t_vol      = 0,
t_store    = 1,    // parent is 1/2 = 0 t_vol
t_page     = 2,    // parent is 2/2 = 1 t_store
t_kv1      = 3,    // parent is 3/2 = 1 t_store
t_record   = 4,    // parent is 4/2 = 2 t_page
t_extent   = 5,
t_user1    = 6,
t_user2    = 7,    // parent is t_user1
t_user3    = 8,    // parent is t_user2
t_user4    = 9     // parent is t_user3
};

struct user1_t {
uint2_t      u1;
    user1_t() : u1(0) {}
    user1_t(uint2_t v1) : u1(v1) {}
};

struct user2_t : public user1_t {
uint4_t      u2;
    user2_t() : u2(0) {}
    user2_t(uint2_t v1, uint4_t v2): user1_t(v1), u2(v2) {}
};

struct user3_t : public user2_t {
uint4_t      u3;
    user3_t() : u3(0) {}
    user3_t(uint2_t v1, uint4_t v2, uint4_t v3)
        : user2_t(v1, v2), u3(v3) {}
};

struct user4_t : public user3_t {
uint4_t      u4;
    user4_t() : u4(0) {}
    user4_t(uint2_t v1, uint4_t v2, uint4_t v3, uint4_t v4)
        : user3_t(v1, v2, v3), u4(v4) {}
};

bool operator==(const lockid_t& p) const;
bool operator!=(const lockid_t& p) const;
friend ostream& operator<<(ostream& o, const lockid_t& i);

uint4_t      hash() const;
void         zero();

name_space_t lspace() const;

```

```

vid_t          vid() const;
const snum_t&   store() const;
const extnum_t& extent() const;
const shpid_t&  page() const;
const slotid_t& slot() const;
uint2_t        u1() const;
uint4_t        u2() const;
uint4_t        u3() const;
uint4_t        u4() const;

void           set_ext_has_page_alloc(bool value);
bool           ext_has_page_alloc() const ;

NORET          lockid_t() ;
NORET          lockid_t(const vid_t& vid);
NORET          lockid_t(const extid_t& extid);
NORET          lockid_t(const stid_t& stid);
NORET          lockid_t(const lpid_t& lpid);
NORET          lockid_t(const stpgid_t& stpgid);
NORET          lockid_t(const rid_t& rid);
NORET          lockid_t(const kvl_t& kvl);
NORET          lockid_t(const lockid_t& i);

NORET          lockid_t(const user1_t& u);
NORET          lockid_t(const user2_t& u);
NORET          lockid_t(const user3_t& u);
NORET          lockid_t(const user4_t& u);

void           extract_extent(extid_t &e) const;
void           extract_stid(stid_t &s) const;
void           extract_lpid(lpid_t &p) const;
void           extract_rid(rid_t &r) const;
void           extract_kvl(kvl_t &k) const;
void           extract_user1(user1_t &u) const;
void           extract_user2(user2_t &u) const;
void           extract_user3(user3_t &u) const;
void           extract_user4(user4_t &u) const;

bool           IsUserLock() const;

void           truncate(name_space_t space);

lockid_t&      operator=(const lockid_t& i);

};

ostream& operator<<(ostream& o, const lockid_t::user1_t& u);
ostream& operator<<(ostream& o, const lockid_t::user2_t& u);
ostream& operator<<(ostream& o, const lockid_t::user3_t& u);
ostream& operator<<(ostream& o, const lockid_t::user4_t& u);

istream& operator>>(istream& o, lockid_t::user1_t& u);
istream& operator>>(istream& o, lockid_t::user2_t& u);
istream& operator>>(istream& o, lockid_t::user3_t& u);

```

```
istream& operator>>(istream& o, lockid_t::user4_t& u);
```

## DESCRIPTION

Locks are acquired implicitly by many **ss\_m** methods. For those situations where more precise control of locking is desired, the following methods allow explicit locking and unlocking.

The class representing a generic lock is a `lockid_t`, described above. The SSM acquires locks on pages, extents, records, stores, and volumes. **The extent locks are NOT to be used by VASs**, simply because the extent-based structure of the SSM is likely to change in future releases.

**lock(lvid, mode, duration, timeout)**

**lock(lockid, mode, duration, timeout)**

The **lock** method is used to acquire a lock on volume, index, file or record. The first version of the method locks the volume specified by *lvid*. The second version locks the index, file or record specified by *lockid*. The *mode* parameter specifies the lock mode to acquire. Valid `lock_mode_t` values are listed in

**basics.h**. The *duration* parameter specifies how long the lock will be held. Valid values (among those listed in **basics.h**) are: `t_instant`, `t_short` and `t_long`. The *timeout* parameter specifies how long to wait for a lock.

**unlock(lockid)**

The **unlock** method releases the most recently acquired lock on the file, index, or record identified by *lockid*. Note, that only locks with duration **t\_short** can be released before end-of-transaction.

**query\_lock(lockid, mode, implicit)**

The **query\_lock** method the mode of the lock held on *lockid* by the current transaction. The lock mode is returned in *mode* and will be **NL** (no lock) if not locked. If *implicit* is **false** then only explicit locks on *lockid* will be considered. For example, if file F is **SH** locked and a query is made about a record in F, the mode returned will be **NL**. **However, if *implicit* is **true**, then **SH** would be returned for this example.**

## Lock Cache Control

Each transaction has a cache of recently acquired locks. The following methods control the use of the cache. These are not supported methods and may be removed in later versions of the software. Note: that the methods only affect the transaction associated with the current thread.

**set\_lock\_cache\_enable(enable)**

The **set\_lock\_cache\_enable** method turns on the cache if *enable* is **true** and turns it off otherwise.

**lock\_cache\_enabled(enabled)**

The **lock\_cache\_enabled** method sets *enabled* to **true** if the lock cache is on.

## Escalation

The lock manager will escalate from a record lock to a page lock, from a page lock to a store lock, and from a store lock to a volume lock, to reduce the number of locks in the table. You can control the thresholds for escalation through the methods **get\_escalation\_thresholds** and **set\_escalation\_thresholds**. The default values are as follows:

record-to-page  
5  
page-to-store  
25  
store-to-volume  
0

In all cases, a threshold of 0 prevents escalation.

When escalation is in use, it be prevented on selected volumes or other lock-able objects through the three **dont\_escalate** methods. If the argument *passOnToDescendants* is *false*, locks acquired on objects below the volume (or given lockid) in the lock hierarchy will still be escalated according to the thresholds.

## ERRORS

TODO

## EXAMPLES

TODO

## VERSION

This manual page applies to Version 2.0 of the Shore Storage Manager.

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## COPYRIGHT

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## SEE ALSO

**transaction(ssm)**, **id(ssm)**, and **intro(ssm)**.