

# Least Squares Cubic Spline Approximation II — Variable Knots

Carl de Boor and John R. Rice

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Department of Computer Sciences  
Purdue University  
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# Least Squares Cubic Spline Approximation II — Variable Knots

Carl de Boor<sup>1</sup> and John R. Rice<sup>2</sup>

## 1 Introduction

This paper presents an extension and application of the algorithm in [2]. We refer to this algorithm as `FIXEDKNOT`. `FIXEDKNOT` is for the computation of least-squares approximations on finite point sets by cubic polynomial splines with fixed knots. The algorithm presented here incorporates `FIXEDKNOT` and treats the knots as variables. The spline depends nonlinearly on the knots and thus we have a nonlinear least-squares approximation problem to solve.

`FIXEDKNOT` has several features specifically designed to facilitate its incorporation into an algorithm which treats the knots as variables. The algorithm presented here is only one of several approaches to varying the knots and some of these other approaches are discussed in general terms at the end of this paper.

An algorithm which varies the knots is desirable because it greatly increases the flexibility of the spline approximants. This flexibility can also be achieved by simply increasing the number of knots. This, of course, increases the complexity of the approximant obtained. There are situations where this increase is an efficient approach since it is difficult and time consuming to solve the nonlinear approximation problem. However, this increase in complexity is very undesirable in an application which involves smoothing or the representation of physical data or shapes. In such an application one must minimize the number of parameters involved in order to obtain the best results. Thus this algorithm is primarily useful for obtaining low to medium accuracy (2 to 5 significant digits) approximations of functions of a more or less arbitrary nature. It has, for example, been used to obtain 3 significant digit approximations to curves with 8 or 10 local extrema and which have a completely unsystematic nature.

Note that this algorithm should *not* be used for high accuracy approximations to mathematically defined functions (e.g., for computer system function subroutines). The degree of convergence for spline approximation is such that this is very unlikely to be efficient unless the polynomial degree is rather high. See [1] and [5] for further results and compare these with polynomial and rational approximations [3].

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## 2 Mathematical Background

We assume that the reader is familiar with FIXEDKNOT and we use the notation of that paper. We recall that a spline of degree  $n$  with  $k$  knots  $\Xi = \{\xi_i | a = \xi_0 < \dots < \xi_{k+1} = b\}$  may be defined by

$$S(A, \Xi, x) = \sum_{i=1}^k a_i (x - \xi_i)_+^n + \sum_{j=0}^n a_{k+j+1} x^j$$

where  $A = (a_1, a_2, \dots, a_{k+n+1})$ .

We consider a function  $f(x)$  defined on a finite set

$$X = \{x_i | a \leq x_i < x_{i+1} \leq b, \quad i = 1, 2, \dots, m\}.$$

Given a value  $n$  for the degree and a number  $h$  of knots we have the

*Approximation Problem.* Determine the spline  $S(A^*, \Xi^*, x)$  so that

$$(2.1) \quad \left[ \int [f(x) - S(A, \Xi, x)]^2 \right]^{\frac{1}{2}}$$

is minimized among all splines of degree  $n$  with  $k$  knots.

Since  $f(x)$  is only defined on the finite set  $X$ , one must use a quadrature formula for the integral in this problem. We assume this is to be done (our algorithm uses the trapezoidal rule), but retain the integral sign for simpler notation.

There are three basic mathematical questions associated with this problem, namely those of the existence, uniqueness and characterization of  $S(A^*, \Xi^*, x)$ . We discuss these briefly.

*The Existence Question.* Simple examples show that this least-squares approximation problem does not always have a solution, e.g., take  $f(x) = |x|$  on  $[-1, +1]$  and approximate by a cubic spline with three knots. One may generalize the concept of spline by allowing the knots to coalesce with the possibility of a resultant loss of smoothness where the knots coalesce. These are called *extended splines* and are presented in [6], see also [4]. In this broader set of approximating functions there always exists a least-squares approximation. In order to avoid technical difficulties, the algorithm presented in this paper does not allow the knots to coalesce.

*The Uniqueness Question.* It is known from general theoretical results [6], from specific theoretical results [C. de Boor, 1963, unpublished] and from examples that the solution of the least-squares nonlinear approximation problem need not be unique. Consider the approximation to  $x^3$  on  $[-1, +1]$  by a broken line with one break. If the break occurs for  $x = 0$ , then by symmetry there is no break. But no least-squares nonlinear spline approximation

can have an inactive knot (see the next section). Thus the best approximation does not have a knot at  $x = 0$  and, again by symmetry, there are at least two best approximations. This line of reasoning can be applied in general.

Furthermore, there may be approximations which are local minima of (2.1), but which are not best approximations. The algorithm presented here attempts to obtain a *local* minimum of (2.1) and hence even if it converges there is no guarantee that a best approximation has been obtained.

*Characterization.* There are no known necessary and sufficient conditions for  $S(A^*, \Xi^*, x)$  to be a best approximation. The algorithm here is based on the usual necessary conditions that one derives for a local minimum.

*Strict Monotonicity of the Error.* It is known [de Boor, 1963, unpublished] for any specific  $f(x)$  and fixed degree  $n$  that if the error

$$E_k^2 = \int [f(x) - S(A^*, \Xi^*, x)]^2$$

of the best approximation with  $k + 1$  knots is not zero, then the error  $E_{k+1}$  of the best approximation with  $k + 1$  knots is strictly less than  $E_k$ , i.e.,  $E_{k+1} < E_k$ . See [4] for details and extensions.

*Inherent Limitations of the Algorithm.* The problem which this algorithm attempts to solve cannot be solved by an algorithm. Thus this algorithm is limited. This theoretical limitation is manifested in several different ways. First, there is the problem of ascertaining when “convergence” has taken place. This is required on two different levels, namely, for the whole algorithm and for the adjustment of knots within this latter problem. The decision that “convergence” has taken place is made on the basis of certain *ad hoc* numerical tests which are not infallible.

These decisions are delicate in view of the need to achieve some efficiency. Thus these tests have been developed on the basis of experience with a certain class of problems. It is hoped that this class is representative of those met in general. However, these tests may be completely inadequate in new situations. If it is intended to use this algorithm extensively for a certain class of problems, it may well pay to experiment with adjustments in these tests in order to achieve better efficiency with minimum risk.

The initial guess for the knots chosen might be extremely poor and result in reaching a local minimum far from the best approximation. The simple scheme of equal spacing used here to obtain an initial guess might well be modified and improved for certain classes of approximation problems.

### 3 The Algorithm and Numerical Procedures

The basic idea of the algorithm is to vary the knots one by one so as to decrease the  $L_2$ -error. This is done systematically from right to left by two procedures, **SWEEP** and **OPT**. The procedure **SWEEP** controls the overall scheme and **OPT** does the variation of the individual knots. The basic scheme used in **OPT** is a discrete Newton's method.

There are two delicate points in an implementation of such a scheme. The first is a suitable choice of termination criteria for the various iterations in the algorithm. One desires to achieve the required accuracy without doing an excessive amount of wasteful computation.

The second point is to make the computation of the  $L_2$ -error as efficient as possible. It is unavoidable that this number be evaluated frequently and it is a nontrivial computation. Furthermore, it is easily seen that it is very inefficient to compute the  $L_2$ -error each time by a standard  $L_2$ -approximation procedure. Note that if only one knot is changed and if only one of the orthogonal functions involves this knot, then the  $L_2$ -approximation problem can be solved on the basis of previous information with an order of magnitude less computation than one can solve such problems in general. It is always arranged so this is the case and the procedure **FIXEDKNOT** has a number of features to allow this. More detailed study shows that it is also possible to make use of some previous information when changing from one knot to another. These points are discussed in more detail in [2].

*Choice of the Initial Knots.* There are two single alternatives. If **NOKNOT** is negative, then  $-\text{NOKNOT}$  knots are chosen equally spaced in the interval  $(\text{XX}(1), \text{XX}(\text{LX}))$ . If **NOKNOT** is positive, then this number of knots is to be read as data. If the function is very unsystematic, it is often profitable to use an initial set of knots concentrated in the regions of rapid change in the function.

*Optimization of the Knots – SWEEP and OPT.* Given an initial set of knots, their optimization is guided by the procedure **SWEEP**. Each knot is, in turn, varied so as to minimize the  $L_2$ -error as a function of this knot. This is started with the last (i.e., the rightmost) interior knot and done sequentially to the left. A cycle refers to one complete pass from right to left. This process is repeated until a termination is encountered.

The variation of the  $I$ -th knot  $\text{XI}(I)$  is carried out in **OPT** using what may be termed the "discrete Newton's" method. Let  $e(t)$  denote the  $L_2$ -error as a function of the position  $t$  of  $\text{XI}(I)$ . Given three points,  $\text{ALEFT} < A < \text{ARIGHT}$ , a new guess **ABEST** for the location of  $\text{XI}(I)$  is determined as the location of the minimum of the parabola  $p(t)$  satisfying

$$p(\text{ALEFT}) = e(\text{ALEFT}), p(A) = e(A), p(\text{ARIGHT}) = e(\text{ARIGHT}).$$

The parabola must have a minimum in order for this to make sense. Also, as to avoid getting wild guesses through extrapolation, **ABEST** should be between **ARIGHT** and **ALEFT**. For this it

is sufficient to have

$$(3.1) \quad e(\text{ARIGHT}), e(\text{ALEFT}) \geq e(\text{A}).$$

Thus the first part of OPT consists of a search algorithm for such a set of three points ARIGHT, ALEFT and A. The basic step size for this search is based on the value of CHANGE = average change in the knots in the preceding cycle. The initial value of CHANGE is .4 and it is measured relative to the length of the interval (XI(I-1), XI(I+1)).

Once such a set is found the parabolic interpolation commences. The newly found guess ABEST replaces one of ARIGHT, ALEFT or A in such a way that the inequalities (3.1) remain valid while making the new value of ARIGHT-ALEFT as small as possible.

*Termination Criteria.* There are two termination criteria for SWEEP. The first is a simple bound on the number of complete cycles or sweeps of varying all the knots, i.e.,

$$(3.2) \quad \text{No more than ITER cycles throught SWEEP}$$

For normal use we recommend that one set ITER=4. In more difficult cases, especially when a larger number of knots is used, one might need to increase ITER. The second criterion is to terminate if

$$(3.3) \quad \left| \frac{\text{PREVER-ERROR}}{\text{ERROR}} \right| \leq .4 * \text{ACC}$$

where ACC = desired accuracy in  $L_2$ -error (*not* the  $L_2$ -error itself), ERROR = current value of the  $L_2$ -error and PREVER = value of the  $L_2$ -error at the start of the current cycle of knot variation. This criterion is based on the assumption that the algorithm is converging linearly (or faster) and the error is reduced at each cycle by a factor of .6 or less. If one notes that the algorithm is converging somewhat slower than this, one should replace the coefficient .4 by a somewhat smaller number.

Note that this is rarely worthwhile to compute an approximation which gives the best  $L_2$ -error with more than one or two significant digits. We recommend setting ACC = .1 for general applications.

There are four termination criterion for OPT. The first is a simple bound on the number of guesses at the best position of XI(I), i.e.,

$$(3.4) \quad \text{No more than INDLP guesses for XI(I)}.$$

We recommend INDLP = 10, a bound which is large enough so that termination rarely occurs from this criterion.

The second criterion is a form of buffering to prevent the knots from coalescing. Set  $H = XI(I+1) - XI(I-1)$ , then constrain  $XI(I)$  by

$$(3.5) \quad XI(I-1) + .0625 H \leq XI(I) \leq XI(I+1) - .0625 H$$

This form of constraint allows a group of knots to become very closely spaced which is sometimes essential. However, it keeps them separated enough to (almost always) avoid failure due to numerical instabilities.

The third criterion is for the search of a triplet of points to initialize the parabolic interpolation phase. The search for such a triplet is terminated if (in the case of search to the right)

$$(3.6) \quad \frac{e(A) - e(ARIGHT)}{ERROR} \leq \frac{ACC}{LXI}$$

where  $LXI$  = number of interior knots and  $ERROR$  is the  $L_2$ -error at the end of the previous cycle. In case of search to the left we terminate if

$$(3.7) \quad \frac{e(A) - e(ALEFT)}{ERROR} \leq \frac{ACC}{LXI}.$$

These criteria are relatively stringent because we feel it is very desirable to be able to enter the parabolic interpolation phase for at least one time. Thus this criterion might not cause termination in `OPT` even when the decrease in the  $L_2$ -error is insignificant for the later stages of the algorithm in a reasonable number of cases.

The criterion can be visualized as based on the assumption that the search is converging linearly (or faster) with an error reduction of  $1-1/LXI$  or smaller. However, the situation here is somewhat different than in `SWEEP` as we do not necessarily desire to expend effort for an accurate placement of  $XI(I)$ . That is to say, in the initial stages of the algorithm the set of knots is far enough from optimum that it is wasteful to accurately optimize one of them with the others inaccurately located. It is unusual for this termination criterion to be active in the terminal phases of the algorithm.

The fourth criterion is for the termination of the parabolic interpolation process. We locate `ABEST` as noted above and compute the value `EPRED` of the parabola at its lowest point, i.e.,  $EPRED = p(ABEST)$ . The optimization is terminated if

$$(3.8) \quad \left| \frac{EPRED - e(ABEST)}{ERROR} \right| \leq 5 * ACC$$

This criterion assumes convergence which is somewhat faster than linear. This is plausible since a discrete Newton method is used. The particular factor 5 was chosen on the basis

of some experiments and reflects a balance between global efficiency and local accuracy as discussed in the preceding paragraph.

The most common cause for termination is that **CHANGE** become small. This implies that little movement of the knots takes place in **OPT** which in turn causes the criterion (3.3) in **SWEEP** to terminate the algorithm.

## 4 Variables in the Program

### Global with **FIXEDKNOT**

|             |             |
|-------------|-------------|
| ADDXI(26)   | LX          |
| COEFL(27,4) | MODE        |
| FCTL(100)   | U(100)      |
| INTERV      | UERROR(100) |
| JADD        | VORDL(28,2) |
| KNOT        | XIL(28)     |
| LMAX        | XX(100)     |
| debug       |             |

### Global in **VARYKNOT**

|        |        |
|--------|--------|
| ACC    | LXI    |
| CHANGE | Q      |
| ERROR  | XI(28) |

### Other Important Variables

|        |          |
|--------|----------|
| A      | INFO(20) |
| ABEST  | INTER    |
| ALEFT  | KVARY    |
| ARIGHT | NOKNOT   |
| EPRED  | PREVER   |
| EPSERR | K        |

### Other Variables

|        |              |
|--------|--------------|
| AA     | ELEFT        |
| AHIGH  | ERIGHT       |
| ALOW   | ETRY         |
| DEL    | II           |
| DELX   | ITRR         |
| DUMB   | K            |
| DXLEFT | LPCNT        |
| DXRGHT | LXI1 = LXI+1 |
| DYLEFT | LXI2 = LXI+2 |
| DYRGHT | SGN          |
| E      | BD           |

## 5 Example

We consider a set of data which has three distinct features: (i) It is actual data (expressing a thermal property of titanium); (ii) It is difficult to approximate using classical techniques;



(iii) There is a significant amount of noise in the data.

**Titanium Heat Data  $XX(I)$ ,  $U(I)$  with approximation  $U^*(I)$   
and error  $UERROR(I)$**

| XX  | U    | U*   | UERROR $\times 10^3$ | XX   | U     | U*    | UERROR $\times 10^3$ |
|-----|------|------|----------------------|------|-------|-------|----------------------|
| 595 | .644 | .619 | 25.35                | 845  | .812  | .796  | 15.86                |
| 605 | .622 | .629 | -7.22                | 855  | .907  | .876  | 31.27                |
| 615 | .638 | .638 | .24                  | 865  | 1.044 | 1.051 | -7.38                |
| 625 | .649 | .644 | 4.50                 | 875  | 1.336 | 1.370 | -34.31               |
| 635 | .652 | .650 | 2.35                 | 885  | 1.881 | 1.838 | 42.64                |
| 645 | .639 | .653 | -14.46               | 895  | 2.169 | 2.195 | -25.98               |
| 655 | .646 | .656 | -10.13               | 905  | 2.075 | 2.078 | -3.19                |
| 665 | .657 | .658 | -.89                 | 915  | 1.598 | 1.582 | 15.62                |
| 675 | .652 | .659 | -6.96                | 925  | 1.211 | 1.197 | 14.15                |
| 685 | .655 | .660 | -4.57                | 935  | .916  | .931  | -14.63               |
| 695 | .664 | .660 | 4.05                 | 945  | .746  | .761  | -15.36               |
| 705 | .663 | .660 | 2.69                 | 955  | .672  | .667  | 5.31                 |
| 715 | .663 | .661 | 2.13                 | 965  | .627  | .624  | 2.71                 |
| 725 | .668 | .662 | 6.12                 | 975  | .615  | .612  | 3.20                 |
| 735 | .676 | .664 | 12.47                | 985  | .607  | .608  | -1.24                |
| 745 | .676 | .666 | 9.93                 | 995  | .606  | .606  | .15                  |
| 755 | .686 | .670 | 16.29                | 1005 | .609  | .604  | 4.62                 |
| 765 | .679 | .675 | 4.32                 | 1015 | .603  | .604  | -.65                 |
| 775 | .678 | .681 | -3.20                | 1025 | .601  | .603  | -2.49                |
| 785 | .683 | .689 | -6.49                | 1035 | .603  | .604  | -.74                 |
| 795 | .694 | .700 | -5.78                | 1045 | .601  | .604  | -3.22                |
| 805 | .699 | .712 | -13.29               | 1055 | .611  | .605  | 6.24                 |
| 815 | .710 | .727 | -17.25               | 1065 | .601  | .605  | -4.19                |
| 825 | .730 | .745 | -14.87               | 1075 | .608  | .605  | 2.66                 |
| 835 | .763 | .765 | -2.39                |      |       |       |                      |

We present two approximations to this data. The first is computed with an initial set of 7 equally spaced knots in the interval (595, 1075). The second is computed with another initial set of knots. This is the approximation shown in the above table.

**Initial Knots**

|                |     |     |     |     |     |      |      |
|----------------|-----|-----|-----|-----|-----|------|------|
| <b>Case 1:</b> | 595 | 675 | 755 | 835 | 915 | 995  | 1075 |
| <b>Case 2:</b> | 595 | 725 | 850 | 910 | 975 | 1040 | 1075 |

The point of these two cases is that the algorithm converges to two distinct local minima of the nonlinear least-squares approximation problem. Note that the data have a very pronounced peak near 900, and in Case 1 we have three interior knots to the left of this peak, while in Case 2 we have only two to the left of this peak.

The final approximations obtained are presented for both cases. The final knots are given along with the coefficients  $\mathbf{C}(I)$ ,  $I = 0, 1, 2, 3$  of the cubic polynomial pieces of the spline. These are the coefficients  $\text{COEFL}(I, J)$ ,  $J = 1, 2, 3, 4$  defined in [2] for the interval  $[\xi_I, \xi_{I+1}]$ . The origin for each polynomial piece is the knot  $\xi_I$ , immediately to the left.

|                    | <b>Case 1</b> |        | <b>Case 2</b>      |   |        |
|--------------------|---------------|--------|--------------------|---|--------|
| Least Square Error | =             | 03489  | Least Square Error | = | .01305 |
| Average Error      | =             | .02296 | Average Error      | = | .00933 |
| Maximum Error      | =             | .11716 | Maximum Error      | = | .04264 |

| KNOTS  | Cubic Coefficients   | KNOTS  | Cubic Coefficients   |
|--------|--|--------|--|
| 595.   | C(0) = .63371<br>C(1) = .16475 <sup>-3</sup><br>C(2) = .19591 <sup>-5</sup><br>C(3) = -.81758 <sup>-8</sup>                | 595.   | C(0) = .61865<br>C(1) = .11658 <sup>-2</sup><br>C(2) = -.11255 <sup>-4</sup><br>C(3) = .37272 <sup>-7</sup>                |
| 755.28 | C(0) = .67678<br>C(1) = .16269 <sup>-3</sup><br>C(2) = -.19723 <sup>-5</sup><br>C(3) = .17472 <sup>-6</sup>                | 835.32 | C(0) = .76609<br>C(1) = .22139 <sup>-2</sup><br>C(2) = .15616 <sup>-4</sup><br>C(3) = .78696 <sup>-5</sup>                 |
| 839.60 | C(0) = .78122<br>C(1) = .35567 <sup>-2</sup><br>C(2) = .42224 <sup>-4</sup><br>C(3) = .92733 <sup>-5</sup>                 | 876.56 | C(0) = .14362 <sup>+1</sup><br>C(1) = .43668 <sup>-1</sup><br>C(2) = .98940 <sup>-3</sup><br>C(3) = -.61055 <sup>-4</sup>  |
| 877.06 | C(0) = .14612 <sup>+1</sup><br>C(1) = .45759 <sup>-1</sup><br>C(2) = .10844 <sup>-2</sup><br>C(3) = -.78416 <sup>-4</sup>  | 902.46 | C(0) = .21703 <sup>+1</sup><br>C(1) = -.27909 <sup>-1</sup><br>C(2) = -.37536 <sup>-2</sup><br>C(3) = .18772 <sup>-3</sup> |
| 896.20 | C(0) = .21844 <sup>+1</sup><br>C(1) = .10478 <sup>-2</sup><br>C(2) = -.34197 <sup>-2</sup><br>C(3) = .91502 <sup>-4</sup>  | 910.47 | C(0) = .18022 <sup>+1</sup><br>C(1) = -.51906 <sup>-1</sup><br>C(2) = .75881 <sup>-3</sup><br>C(3) = -.37241 <sup>-5</sup> |
| 910.22 | C(0) = .17793 <sup>+1</sup><br>C(1) = -.40887 <sup>-1</sup><br>C(2) = .42760 <sup>-3</sup><br>C(3) = -.13771 <sup>-5</sup> | 977.85 | C(0) = .61061<br>C(1) = -.37235 <sup>-3</sup><br>C(2) = .60407 <sup>-5</sup><br>C(3) = -.28471 <sup>-7</sup>               |
| 1075.  |  | 1075.  |  |

The algorithm required six cycles through SWEEP for Case 1. The  $L_2$ -error decreased as follows:

| cycle        | 1      | 2      | 3      | 4      | 5      | 6      |
|--------------|--------|--------|--------|--------|--------|--------|
| $L_2$ -error | .09176 | .05927 | .03944 | .03588 | .03509 | .03489 |

The algorithm required seven cycles through SWEEP for Case 2. The  $L_2$ -error decreased as follows:

|              |        |        |        |        |        |        |        |
|--------------|--------|--------|--------|--------|--------|--------|--------|
| cycle        | 1      | 2      | 3      | 4      | 5      | 6      | 7      |
| $L_2$ -error | .04595 | .03848 | .02761 | .02177 | .01432 | .01321 | .01305 |

These two cases required about 17 and 23 seconds, respectively, of execution time on a IBM 7094 for a FORTRAN IV version of this algorithm. They required about xxx and yyy seconds, respectively, of execution time on a CDC 6500 in Algol.

## 6 Other Nonlinear Algorithms Based on FIXEDKNOT

The procedure FIXEDKNOT is designed to be readily adaptable to form a basis for a variety of nonlinear spline approximation algorithms. We briefly outline four such algorithms. We have used one of these (the last one) extensively and another (the second one) in some experimentations.

### 6.1 Non-systematic knot optimization

We have observed that there is frequently a significant amount of wasted computation in problems involving a larger number of knots, say more than 5 or 6. It occurs that a few, perhaps most, of the knots become correctly placed, while the remaining ones (somewhat more delicate) requires several additional cycles to locate accurately. The systematic nature of the algorithm VARYKNOT requires one nevertheless to adjust the position of all knots in each cycle. It is clear to us that one can devise workable criteria for determining reasonably well which knots are more critical. One could use these criteria to optimize the knots in an unsystematic manner to increase the efficiency of the computation. We have not formalized such criteria and believe their use would significantly increase the logical complexity of the algorithm.

### 6.2 Systematic insertion of additional knots – $L_\infty$ criterion

A plausible scheme is to start out with no knots at all, find the best linear cubic approximation, then insert a knot near (or at) the location of the maximum error. One then could compute a linear spline approximation with one knot, and insert a second knot near the location of the maximum error. This process is then repeated until the error is reduced to some desired level.

We have experimented with this scheme and it does in fact work. Special provisions must be made if the data contain wild points or pronounced peaks. The maximum error will then

occur several times at one point. The new knots should be placed on alternating sides of this point and prevented from converging to this point. It usually happens that enough knots are placed in the neighborhood of a wild point so that the data are actually interpolated nearby. This is normally undesirable and this scheme is not recommended for such data.

This scheme is not as attractive as we had expected. In addition to the problem of wild points and peaks, it consistently leads to more knots than really required, sometimes excessively so. However, it usually requires less computation time than schemes (e.g., see 6.4) which optimize the locations of knots. Thus when this process was applied to the data of the example, it took 15 interior knots to produce an approximation of the same accuracy as had been obtained in Case 2 above with an optimal placing of 5 interior knots. Execution time, on the other hand, on an IBM 7094, was merely 4 seconds. We conclude that the location of the maximum error is not a completely reliable guide for the place to insert additional knots.

### **6.3 Systematic insertion of additional knots – $L_2$ criterion**

Consider a process like 6.2 where we locate that interval between adjacent knots which has the most error in the  $L_2$  sense. We suspect that it is better to insert additional knots into this interval than near the location of the maximum error. We have not tested this suspicion, however.

### **6.4 Systematic insertion of knots with optimization**

We have used extensively an algorithm which systematically increases the number of knots and optimizes all knots after each insertion. This algorithm only requires the user to specify the desired accuracy of approximation and the algorithm determines the number as well as the location of the knots. In order to achieve efficiency, the convergence criteria during the algorithm must depend on how close one is to the requested accuracy. Once this matter is satisfactorily settled, we find that it requires only slightly longer to obtain suitable approximations with this scheme than it does with VARYKNOT starting with the correct number of knots roughly placed.

Note that the algorithm is essentially different from that of 6.2. Even though the initial guess at the new knot locations is made similarly, the optimization process eliminates the difficulties with wild points. In fact, it is highly recommended for data smoothing, the identification of wild points and other types of data analysis.

## 7 References

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C PROGRAM SPLINE(INPUT,OUTPUT,TAPE5=INPUT,TAPE6=OUTPUT)
C NONLINEAR SPLINE APPROXIMATION
C PROGRAM WRITTEN BY CARL DE BOOR AND JOHN RICE
C PURDUE UNIVERSITY
C SUPPORTED BY THE NATIONAL SCIENCE FOUNDATION GP-4052,GP-7163
C
C PLEASE REPORT ANY CASES OF INOPERATION TO THE AUTHORS.
C THANKS
C **** NUMERICAL ANALYSIS CONTROL ****
C CONTROL PARAMETERS FUNCTION
C ITER NO. OF SWEEPS THRU OPT
C BD (IN OPT) IMPROVEMENT NEEDED TO REPEAT
C EPSERR( IN SWEEP) IMPROVEMENT NEEDED TO REPEAT
C DIST (IN OPT NEAR 30,80) KEEPS KNOTS SEPARATED
C INDLP NO. OF PASSES THRU OPT
C THE FOLLOWING IS THE MAIN PROGRAM FOR VARYKNOT
C
C IMPLICIT NONE
C CHARACTER *4 INFO(20)
C LOGICAL debug
C real ACC,ADDXI(26),CHANGE,COEFL(27,4),DEL,DUMB,ERROR,
1 FCTL(100),
2 FXDKNT,
3 U(100),UERROR(100),VORDL(28,2),Q,
4 XI(28),XIL(28),XX(100)
C integer I,IABS,IERROR,INTERV,ITER,
1 J,JADD,KNOT,L,LMAX,LX,LXI,LXI1,LXI2,MODE,NOKNOT
C
C COMMON INPUT SERVES AS INPUT TO FXDKNT
C SEE FXDKNT FOR DEFINITIONS OF VARIABLES
C COMMON/INPUT/LX,XX,U,JADD,ADDXI,MODE,debug
C COMMON OUTPUT SERVES AS OUTPUT FROM FXDKNT
C SEE FXDKNT FOR DEFINITIONS OF VARIABLES
C COMMON/ OUTPUT /UERROR,FCTL,XIL,COEFL,VORDL,KNOT,LMAX,INTERV
C
C COMMON OTHER SERVES AS COMMUNICATION BETWEEN OPT,SWEEP AND HERE
C LXI = NUMBER OF INTERIOR KNOTS, LXI1 = LXI+1, LXI2 = LXI+2

```

```

C      Q  = NUMERICAL CONTROL VARIABLE USED BETWEEN OPT AND SWEEP
C      CHANGE = DITTO
C      ERROR = CURRENT VALUE OF THE L-2 ERROR - SQUARED
C      ACC  = DESIRED ACCURACY OF L-2 ERROR
C      XI(28)= ARRAY FOR KNOTS
COMMON/ OTHER / LXI,LXI1,LXI2,Q ,CHANGE,ERROR ,ACC, XI
debug = .false.

C
C      ACC = .1 AND ITER = 4 TO 8 SEEM TO BE GOOD VALUES FOR TYPICAL USES
      ACC = .1
      ITER = 8
      ITER = 2      ! temporary

C
C ***INFO IS SIMPLY AN IDENTIFICATION OF THE DATA***
      1 READ(5,550, END = 40), (INFO(I), I=1,20)
550 FORMAT(20A4)
      WRITE(6,551), (INFO(I), I=1,20)
551 FORMAT( 20A4)

C
C      READ IN NO. OF POINTS=LX AND THE DATA XX AND U
C *** IF NOKNOT.GE.2, THEN READ IN LXI2=NOKNOT KNOTS***
C *** OTHERWISE PROGRAM CHOOSES LXI2 =-NOKNOT EQUISPACED KNOTS ***
      READ *, NOKNOT, LX, (XX(I), U(I), I=1,LX)
      LXI2 = IABS(NOKNOT)
C IF *NOKNOT* IS .GT. 0, READ IN NOKNOT KNOTS (INCL.BOUNDARY POINTS.)
      IF (NOKNOT .GT. 0) READ *, (XI(J),J=1,LXI2)

C
C      **CHECK ON GIVEN DATA
C      THESE CHECKS PREVENT USER FROM EXCEEDING BOUNDS ON STORAGE
C      AND FROM PRESENTING UNORDERED XX ARRAY
      IERROR = 0
      IF (LX .GE. LXI2+2 .AND. LX .LE. 100) GO TO 3
      WRITE(6,662) LX
662 FORMAT(' NO. OF DATA POINTS, LX = ',I4,
      1     ',NOT WITHIN THE BOUNDS ABS(NOKNOT)+2 AND 100')
      IERROR = 1

      GO TO 7

```



```

3 IF (LXI2 .GE. 3 .AND. LXI2 .LE. 28) GO TO 4
  WRITE(6,660) NOKNOT
660 FORMAT('1KNOT CONTROL PARAMETER NOKNOT =',I3,
1      ' NOT WITHIN BOUNDS.')
```

IERROR = 1

```

4 DO 6 L=2,LX
  IF (XX(L) .GT. XX(L-1))      GO TO 6
  WRITE(6,664) L,XX(L),U(L)
664  FORMAT(' DATA POINT ',I4,2F14.8,' NOT IN ASCENDING ORDER.')
```

IERROR = IERROR + 1

```

6  CONTINUE
  IF (IERROR .LT. 1)          GO TO 14
7  WRITE(6,666) IERROR
666 FORMAT(' *** CORRECT INDICATED ',I3,
1      ' INPUT ERROR(S) AND RESTART.')
```

GO TO 1

C

C

\*\*INITIALIZE

```

14 IF (NOKNOT .GT. 0)          GO TO 30
```

C

C

WHEN NOKNOT IS NEG., INTRODUCE -NOKNOT EQUISPACED KNOTS

XI(1) = XX(1)

XI(LXI2) = XX(LX)

DEL = (XX(LX) - XX(1))/FLOAT(LXI2-1)

DO 26 J = 3,LXI2

XI(J-1) = XI(J-2) + DEL

26 continue

C

C

SET UP INITIAL APPROXIMATION

```

30 ADDXI(1) = XI(1)
```

ADDXI(2) = XI(LXI2)

LXI1 = LXI2-1

LXI = LXI1-1

MODE = 0

JADD = LXI2

DO 35 J = 3,LXI2

ADDXI(J) = XI(J-1)

```

35 continue
   ERROR = FXDKNT(0.)
C   ***NOTE.  MODE HAS BEEN SET EQUAL TO 1
C   *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
   WRITE(6,611) (L,XX(L),U(L),UERROR(L), L=1,LX)
611 FORMAT(' GIVEN DATA AND ERROR IN FIRST APPROXIMATION'//
1      (I4,3F14.8))
   WRITE(6,612) NOKNOT,ITER
612 FORMAT(' '//' NO. OF INITIAL KNOTS =',I3/
1      ' ITER =',I3)
   WRITE(6,900) (XI(I), I=1,LXI2)
900 FORMAT(' KNOTS PRIOR TO OPTIMIZATION'/(9F12.6))
C
C           OPTIMIZE KNOTS
   CALL SWEEP(ITER)
C
   WRITE(6,640)
640 FORMAT(49X,'*** FINAL OUTPUT ***'////)
   MODE = 1
   JADD = 0
   DUMB = FXDKNT(1.)
                                     GO TO 1
40                                     STOP
C
   END
C
C*****
C
   SUBROUTINE SWEEP(ITRR)
C
C           KVARY+1 = INDEX OF KNOT           BEING VARIED
C           SUBROUTINE OPT(I) OPTIMIZES ITH INTERIOR KNOT
C
   IMPLICIT NONE
   LOGICAL debug
   real ACC,ADDXI(26),CHANGE,COEFL(27,4),DUMB,EPSERR,ERROR,
1      FCTL(100),

```

```

2     FXDKNT,PREVER,Q,
3     U(100),UERROR(100),VORDL(28,2),
4     XI(28),XIL(28),XX(100)
      integer I,INTERV,ITER,ITRR,
1     JADD,K,KNOT,KVARY,LMAX,LX,LXI,LXI1,LXI2,MODE
      COMMON/INPUT/LX,XX,U,JADD,ADDXI,MODE,debug
      COMMON/ OUTPUT /UERROR,FCTL,XIL,COEFL,VORDL,KNOT,LMAX,INTERV
      COMMON/ OTHER / LXI,LXI1,LXI2,Q ,CHANGE,ERROR ,ACC,      XI
C     AT ALL TIMES, ERROR CONTAINS (L2 ERROR)**2 OF CURRENT B.A.
C
      ITER = ITRR
C     **NEXT      CARDS SET NUMERICAL ANALYSIS CONTROLS
      EPSERR =      ACC/2.5
      CHANGE = .4*FLOAT(LXI)
C
10  KVARY = LXI
      Q = CHANGE/FLOAT(LXI)
      if (.not. debug)      go to 11
C     *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
      WRITE (6,902) ITER,Q
902  FORMAT (' ITER, Q ',I5,E20.8)
11  CHANGE = 0.
      PREVER = ERROR
      MODE = 2
      JADD = 0
      KNOT = KNOT - 1
      DUMB = FXDKNT(0.)
20  CONTINUE
      if (.not. debug)      go to 21
C     *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
      WRITE(6,900) KVARY
900  FORMAT(' '///' VARYING',I4,'TH INTERIOR KNOT')
      WRITE(6,901) ERROR
901  FORMAT(' SQ. OF L2-ERROR ',E16.6)
C
21  CALL OPT(KVARY)
      KVARY = KVARY -1

```

```

JADD = JADD + 1
IF( JADD .LE. 1 )      GO TO 25
K= JADD
DO 22 I = 2,JADD
K= K-1
22 ADDXI(K+1) = ADDXI(K)
25 ADDXI(1) = XI(KVARY + 2)
KNOT = LXI1 - JADD
MODE = 2
DUMB = FXDKNT(0.)
IF( KVARY .NE. 0 )      GO TO 20
C      THE LAST CALL TO FXDKNT PRODUCES THE B.A. USING ALL KNOTS,
C      SINCE THEN ADDXI CONTAINS ALL KNOTS
ERROR = DUMB
C      *** THE FOLLOWING TWO CARDS PRODUCE PRINTED OUTPUT OF L1,L2,L-INF
C *      JADD = 0
C *      DUMM = FXDKNT(2.)
C
C      **IF CHANGE IN ERROR IS BIG ENOUGH MAKE ANOTHER SWEEP, ELSE QUIT
IF (PREVER-ERROR .LE. EPSERR*PREVER)      GO TO 60
ITER = ITER-1
C
C      **CHECK NUMBER OF PASSES THROUGH SWEEP
IF(ITER.GT.0)      GO TO 10
40 CONTINUE
C
C      IN FINAL VERSION GO TO 40, GO TO 60 ARE REPLACED BY RETURN
if (.not. debug)      return
C      *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
WRITE(6,620)
620 FORMAT (' *** NO. OF ALLOWABLE SWEEPS USED UP')
      RETURN
60 CONTINUE
if (.not. debug)      return
C      *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
WRITE(6,610)
      RETURN

```

```

610 FORMAT(' *** SWEEP DISCONTINUED - INSUFFICIENT CHANGE IN ERROR')
      END
C
C*****
C
      SUBROUTINE OPT(II)
C
C      I REFERS TO THE ITH INTERIOR KNOT
C      OPT FINDS THE OPTIMAL ITH KNOT BETWEEN THE I-1ST AND I+1ST KNOTS
C      THE REMAINING KNOTS ARE HELD FIXED.
C      INDLP = A BOUND ON THE NUMBER OF TRIES ALLOWED
C              FOR IMPROVEMENT OF THE ITH KNOT
C      Q = MULTIPLICATION FACTOR WHICH SHOULD DECREASE AS A
C          FUNCTION OF THE NO. OF SWEEPS THRU SWEEP
C          Q IS ALTERED IN SWEEP
C
      IMPLICIT NONE
      LOGICAL debug
      real A,AA,ABEST,ACC,ADDXI(26),AHIGH,ALEFT,ALOW,ARIGHT,BD,
1      CHANGE,COEFL(27,4),DEL,DIFF,DIST,DXLEFT,DXRGHT,DYLEFT,
2      DYRGHT,E,EBEST,ELEFT,EPRED,ERIGHT,ERROR,ETRY,FCTL(100),
3      FXDKNT,H,Q,SGN,U(100),UERROR(100),VORDL(28,2),
4      XI(28),XIL(28),XX(100)
      integer I,II,INDLP,INTERV,
1      JADD,KNOT,LMAX,LPCNT,LX,LXI,LXI1,LXI2,MODE
      COMMON/INPUT/LX,XX,U,JADD,ADDXI,MODE,debug
      COMMON/ OUTPUT /UERROR,FCTL,XIL,COEFL,VORDL,KNOT,LMAX,INTERV
      COMMON/ OTHER / LXI,LXI1,LXI2,Q,CHANGE,ERROR,ACC,XI
C
      I = II
C      **NUMERICAL ANALYSIS PARAMETERS SET HERE
      INDLP=9
      BD = ACC*ERROR/FLOAT(LXI)
      DIST = .0625
      H = XI(I+2)-XI(I)
      ALOW = XI(I) + DIST*H
      AHIGH = XI(I+2) - DIST*H

```

```

LPCNT= 0
MODE = 3
C
C   **BEGIN SEARCH - FIND THREE VALUES FOR THE ITH KNOT
C   SUCH THAT L2-ERROR AT MIDDLE VALUE, A , IS LESS THAN
C   ERROR AT LEFT VALUE, ALEFT, AND AT RIGHT VALUE, ARIGHT
A = XI(I+1)
E = FXDKNT(A)
ALEFT = A + Q*(XI(I)-A)
ELEFT = FXDKNT(ALEFT)
if (.not. debug)          go to 5
C   *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
ARIGHT = 0.
ERIGHT = 0.
WRITE (6,900) ELEFT,E,ERIGHT,ALEFT,A,ARIGHT
5 SGN = SIGN(1.,ELEFT-E)
IF (SGN.GE.0.)           GO TO 20
                        GO TO 60

C
C   **SEARCHING FOR NEW KNOT TO THE RIGHT
10 ALEFT = A
ELEFT = E
A = ARIGHT
E = ERIGHT
20 ARIGHT = A + Q*(XI(I+2)-A)

C
C   **BUFFER TO PREVENT COALESCING OF KNOTS
30 IF (AHIGH.GE.ARIGHT)  GO TO 40
AA = AHIGH
if (.not. debug)        go to 199
C   *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
WRITE(6,610) I
                        GO TO 199

C
40 ERIGHT = FXDKNT(ARIGHT)
if (.not. debug)        go to 41
C   *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***

```

```

        WRITE (6,900) ELEFT,E,ERIGHT,ALEFT,A,ARIGHT
41 IF (E.LE.ERIGHT)          GO TO 100
C
C      **CHECK TO STOP OPT
      IF(E -ERIGHT.LE.BD .OR. LPCNT .GT. INDLP ) GO TO 240
50 LPCNT = LPCNT+1
      IF(SGN.GT.0.) GO TO 10
C
C      **SEARCHING FOR NEW KNOT TO THE LEFT
60 ARIGHT = A
      ERIGHT = E
      A = ALEFT
      E = ELEFT
70 ALEFT = A + Q*(XI(I)-A)
C
C
C      **BUFFER TO PREVENT COALESCING OF KNOTS
80 IF (ALEFT.GE.ALLOW)          GO TO 90
      AA = ALLOW
      if (.not. debug)          go to 199
C      *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
      WRITE(6,620) I
                                   GO TO 199
C
90 ELEFT = FXDKNT(ALEFT)
      if (.not. debug)          go to 91
C      *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
      WRITE (6,900) ELEFT,E,ERIGHT,ALEFT,A,ARIGHT
91 IF (E.LE.ELEFT)          GO TO 100
C
C      **CHECK TO STOP OPT
      IF(E - ELEFT.LE.BD .OR. LPCNT .GT. INDLP ) GO TO 230
                                   GO TO 50
C
C      **REQUIRED 3 VALUES HAVE BEEN FOUND
C      FOLLOWING CODE FINDS PT. AT WHICH MIN OF PARABOLA CURVE PASSIN
C      THRU THE ERROR VALUES AT THE PTS ALEFT, A, ARIGHT OCCURS

```

```

100 DXLEFT = ALEFT - A
    DXRGHT = ARIGHT - A
    DYLEFT = (ELEFT-E)/DXLEFT
    DYRGHT = (ERIGHT-E)/DXRGHT
    DIFF = DYLEFT - DYRGHT
    IF (DIFF .EQ. 0.)                GO TO 200
    DEL = .5/DIFF*(DXRGHT*DYLEFT-DXLEFT*DYRGHT)
    EPRED = E+DEL*(DYRGHT+(DXRGHT-DEL)/(ARIGHT-ALEFT)*DIFF)
    ABEST = A + DEL
    EBEST = FXDKNT(ABEST)
    if (.not. debug)                go to 109
C   *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
    WRITE (6,900) ELEFT,EBEST,ERIGHT,ALEFT,ABEST,ARIGHT
C
C   **DETERMINE WHETHER ABEST GIVES BEST APPRX AND MAKE APPROPRIATE
C   SWITCHING OF THE AI'S DEPENDING ON SIGN OF DEL
109 IF (EBEST.LE.E)                GO TO 130
    IF(DEL)110,200,120
110 ALEFT = ABEST
    ELEFT = EBEST
    GO TO 170
120 ARIGHT = ABEST
    ERIGHT = EBEST
    GO TO 170
130 IF(DEL)140,200,150
140 ARIGHT = A
    ERIGHT = E
    GO TO 160
150 ALEFT = A
    ELEFT = E
160 A = ABEST
    E = EBEST
C
C   **FOLLOWING          TESTS DETERMINE WHETHER OR NOT TO
C   REITERATE PARABOLA MINIMIZATION PHASE
170 IF (ABS(EPRED-EBEST).LT.5.*BD)  GO TO 210
    IF(LPCNT.GT.INDLP)              GO TO 200

```



```

        LPCNT = LPCNT+1
                                GO TO 100
C
199 ETRY = FXDKNT(AA)
    IF (E.LT.ETRY)              GO TO 200
    A = AA
    E = ETRY
200 CHANGE = CHANGE + ABS(A -XI(I+1))/H
    XI(I+1) = A
    ERROR = E
    if (.not. debug)            RETURN
C    *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
    WRITE (6,900) ELEFT,E,ERIGHT,ALEFT,A,ARIGHT
                                RETURN
C
C        IN FINAL VERSION GO TO 210, IS REPLACED BY GO TO 200
210 CONTINUE
    if (.not. debug)            go to 200
C    *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
    WRITE(6,640) LPCNT
    GO TO 200
230 A = ALEFT
    E = ELEFT
    if (.not. debug)            go to 200
C    *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
    WRITE(6,640) LPCNT
    GO TO 200
240 A = ARIGHT
    E = ERIGHT
    if (.not. debug)            go to 200
C    *** THIS IS TEMPORARY DEBUGGING AND TESTING OUTPUT ***
    WRITE(6,640) LPCNT
    GO TO 200
610 FORMAT(' *** OPT DISCONTINUED - KNOT BEING OPTIMIZED (',I2,
1') MOVED TOO CLOSE TO RIGHT NEIGHBOR')
620 FORMAT(' *** OPT DISCONTINUED - KNOT BEING OPTIMIZED (',I2,
1') MOVED TOO CLOSE TO LEFT NEIGHBOR')

```

```
640 FORMAT(' *** OPT DISCONTINUED AT',I4,  
1' - INSUFFICIENT CHANGE IN ERROR')  
900 FORMAT(' PARABOLA - ERROR VALUES ',3E20.6/  
112X,'AI VALUES      ',3E20.6)  
END
```

```

C
C*****
C*****
C*****
C
      FUNCTION FXDKNT (CHANGE)
C           THE FUNCTION RETURNS THE SQUARE OF THE L2-ERROR
      IMPLICIT NONE
      LOGICAL MODE3
C** IT MAY BE NECESSARY ON SOME SYSTEMS TO MENTION ALL COMMON BLOCKS
C LISTED HERE IN THE PROGRAM CALLING *FXDKNT*, TOO, TO INSURE THAT
C THE INFO IN THESE BLOCKS DOES NOT DIE BETWEEN CALLS TO *FXDKNT*.
      real ADDXI(26),BC(30),CHANGE,COEFL(27,4),CUBERR(100),DIF,
1      DOT,ERBUT1,ERRL1,ERRL2,ERRL99,FCT(100,30),FCTL(100),
2      FXDKNT,SCALE,SERROR(100),T,TREND(100),
3      U(100),UERROR(100),VORD(30,28,2),VORDL(28,2),W,
4      WEIGHT(100),XIL(28),XKNOT,XSCALE,XX(100)
      integer I,IDUM,IE,I0,ILAST,ILM3,ILOC,INSERT,INSIRT(30),INTERV,
1      IORDER(28),IPRINT,J,JADD,K,KNOT,KNOTSV,L,LMAX,LX,MODE
      DOUBLE PRECISION TRPZWT(100),SUM
      COMMON / WANDT / TREND,TRPZWT
      COMMON/ INPUT /LX,XX,U,JADD,ADDXI,MODE
C           U(L) = FCT TO BE APPR AT XX(L), L=1,LX.
C           XX(L) IS ASSUMED TO BE NONDECREASING WITH L
C           ADDXI(I) = I-TH KNOT TO BE INTRODUCED, I=1,JADD
C           MODE = 0,1,2,3 . SEE COMMENTS BELOW ( AND IN NUBAS)
      COMMON/ OUTPUT /UERROR,FCTL,XIL,COEFL,VORDL,KNOT,LMAX,INTERV
C           UERROR(L) = ERROR OF B.L2 A. TO U, L=1,LX
C           KNOT = CURRENT NO. OF KNOTS (INCL BDRY KNOTS)
C           INTERV = KNOT - 1 = CURRENT NO. OF INTERVALS (POL.PIECES)
C           XIL(K),K=1,KNOT, CURRENT (ORDERED) SET OF KNOTS
C           THE MAXIMUM ERROR OCCURS AT XX(LMAX)
C           IF ARG=1, FCTL(L) CONTAINS THE CURRENT B.A TO U AT XX(L)
C           COEFL(I,.) CONTAINS THE POL.COEFL. ON I-TH INTERVAL FOR B.A.
C           VORDL(I,.) CONTAINS VALUE AND DERIV. OF B.A. AT XIL(I)
      COMMON/ BASIS /FCT,VORD,BC,ILAST
C** A CHANGE IN THE COLUMN LENGTH OF *FCT* FORCES CHANGE IN ST.NO.69

```

```

C   IN *NUBAS* .
C       FCT (L,M) = BASIS FCT M AT XX(L)
C       VORD(M,K,L) CONTAINS THE ORDS (L=1) AND SLOPES (L=2) OF FCT M
C       AT THE KNOT INTRODUCED AS K-TH. CORRELATION TO ORDERING OF
C       KNOTS BY SIZE IS DONE VIA IORDER, I.E., ORD AND SLOPE AT
C       XIL(K) ARE IN VORD(M,IORDER(K),.).
C       BC(I) = COORDINATE OF U (AND OF B.A. TO U) WRTO I-TH O.N.FCT
C       ILAST = CURRENT NO. OF BASIS FCTNS
COMMON/ LASTB /IORDER,INSIRT,XKNOT
C       THE FCT ILAST (TO BE) INTRODUCED LAST HAS ADDITIONAL KNOT
C       XKNOT, THE KNOT JUST INTRO-
C       DUCED HAS INDEX INSERT IN XIL,INSERT IS SAVED IN INSIRT(ILAS
C       FOR POSSIBLE REPLACEMENT OF KNOTS LATER ON (SEE MODE=2,3).
C   ***LOCAL VARIABLES
COMMON /LOCAL/ XSCALE,KNOTSV,ERBUT1,CUBERR,WEIGHT,MODE3
C       XSCALE = XX(LX) - XX(1), USED TO NORMALIZE INNER PRODUCT
C       = LENGTH OF THE INTERVAL OF INTEGRATION
C       KNOTSV = NO. OF KNOTS USED IN MOST RECENT CALL TO FXDKNT
C       ERBUT1 = SQ. OF L2-ERROR OF APPR USING ALL BUT THE ONE
C       KNOT BEING VARIED ( USED IN MODE = 3)
C       CUBERR = UERROR OF B.A. BY CUBIC POL-S (NEEDED FOR MODE = 2)
C       MODE3 = TRUE OR FALSE DEP. ON WHETHER PREV. CALL WAS IN
C       MODE=3 OR NOT
C       CHANGE = THE NEW VALUE OF THE KNOT BEING VARIED IF MODE=3;
C       IT IS USED TO CONTROL PRINTED OUTPUT OTHERWISE.
IPRINT = IFIX(CHANGE)
IF (MODE.GT.0)                GO TO 29
C-----
C       *** MODE=0* COMPUTE BASIS FCTNS 1 THROUGH 4 AND B.A. TO U WRTO
C       THESE, THEN SET MODE = 1 AND PUT UERROR INTO U.
XSCALE = XX(LX) - XX(1)
DO 10 I=5,30
    INSIRT(I) = 0
10 continue
DO 11 L=1,LX
    UERROR(L) = U(L)
    TREND(L) = T(XX(L))

```

```

        WEIGHT(L) = W(XX(L))
11 continue
    DO 12 L=3,LX
        TRPZWT(L-1) = (XX(L)-XX(L-2))/2.*WEIGHT(L-1)
12 continue
    TRPZWT(1) = (XX(2)-XX(1))/2.*WEIGHT(1)
    TRPZWT(LX) = (XX(LX)-XX(LX-1))/2.*WEIGHT(LX)
C
    XIL(1) = ADDXI(1)
    XIL(2) = ADDXI(2)
    IORDER(1) = 1
    IORDER(2) = 2
    KNOT = 2
    INTERV = 1
    DO 19 I=1,4
        ILAST = I
        CALL NUBAS
        DO 19 L=1,LX
            UERROR(L) = UERROR(L) - BC(I)*FCT(L,I)
19 continue
C
    MODE = 1
    DO 20 L = 1,LX
        CUBERR(L) = UERROR(L)
20 continue
C    IF (JADD.LE.2), ONLY B.A. BY CUBICS IS COMPUTED
C    OTHERWISE, ADDXI(I), I.GT.2, CONTAINS ADDITIONAL KNOTS
    JADD = JADD - 2
    IF (JADD.LE.0)                GO TO 60
    DO 21 I=1,JADD
        ADDXI(I) = ADDXI(I+2)
21 continue
                                                GO TO 51
C-----
    29                                GO TO (40,40,30),MODE
C-----
C    *** MODE=3 *** MERELY CHANGE THE LAST KNOT INTRODUCED TO

```

```

C          CHANGE (THE INPUT ARGUMENT TO  FIXDKNT ) AND
C          RECOMPUTE L2 ERROR.
C          THIS MODE SHOULD BE USED FOR MINIMIZING THE
C          L2-ERROR WRTO THE KNOT
C          INTRODUCED LAST AS IT MINIMIZES THE COMP WORK
C          IF MODE3 = TRUE (I.E., THE PRECEDING CALL TO FXDKNT
C          WAS IN MODE=3),THE PROGR WILL ASSUME THAT CHANGE
C          HAS THE SAME ORDER REL TO THE OTHER KNOTS AS THE
C          PREV.INTRODUCED VALUE FOR KNOT. OTHERWISE
C          IF MODE3 = FALSE, I.E., THE PRECEDING CALL WAS IN
C          SOME OTHER MODE), A FCT IS ADDED WITH CHANGE AS
C          THE ADDITIONAL KNOT.
C          UERROR IS ASSUMED TO CONTAIN ERROR OF B.A. TO U
C          ALL PREV FCTNS, AND ERBUT1 the SQ. OF ITS L2NORM
C          **NOTE** IF THE NEXT CALL TO FXDKNT
C          IS IN A MODE OTHER THAN 3, THE CHANGE PROPOSED
C          NOW WILL BE MADE PERMANENT.

```

```

30 XKNOT = CHANGE
    IF (MODE3)
        MODE3 = .TRUE.
        MODE = 2
        CALL NUBAS
        KNOTSV = KNOT
        MODE = 3

```

```

GO TO 35

```

```

GO TO 36

```

```

35 CALL NUBAS

```

```

36 FXDKNT = ABS(ERBUT1 - BC(ILAST)/XSCALE*BC(ILAST))

```

```

RETURN

```

```

C-----

```

```

C          ***MODE=1,2***  RETAIN THE FIRST KNOT KNOTS INTRODUCED EARLIER
C          (HENCE THEIR CORRESP FCTNS) BUT REPLACE FURTHER
C          FCTNS (IF ANY) BY FCTNS HAVING ADDITIONAL
C          KNOTS ADDXI(I),I=1,JADD) HENCE
C          IF KNOT.LT.KNOTSV(=NO.OF KNOTS USED IN PREV CALL
C          THEN 40 THROUGH 49 RESTORES ARRAYS IORDER,XIL,
C          UERROR TO THE STATE OF  ILAST = KNOT + 2 ,
C          INVERTING THE ACTION OF DO 11 ... TO 14 IN NUBAS

```

```

40 IF (KNOT.LT.KNOTSV)                GO TO 42
    KNOT = KNOTSV
    IF (.NOT.MODE3)                    GO TO 50
    DO 41 L=1,LX
        UERROR(L) = UERROR(L) - BC(ILAST)*FCT(L,ILAST)
41 continue
                                                    GO TO 49
42 DO 43 L=1,LX
        UERROR(L) = CUBERR(L)
43 continue
    IF (KNOT.LE.2)                    GO TO 48
    IDUM = KNOT + 1
    DO 45 IO=IDUM,KNOTSV
        INSERT = INSIRT(ILAST)
        ILM3 = ILAST - 3
        DO 44 K=INSERT,ILM3
            IORDER(K) = IORDER(K+1)
            XIL(K) = XIL(K+1)
44 continue
        ILAST = ILAST-1
45 continue
    DO 47 I=5,ILAST
        DO 47 L=1,LX
            UERROR(L) = UERROR(L) - BC(I)*FCT(L,I)
47 continue
                                                    GO TO 49
48 XIL(2) = XIL(ILAST-2)
    IORDER(2) = 2
    KNOT = 2
49 IF (JADD.GT.0)                    GO TO 51
    ILAST = KNOT + 2
    INTERV = KNOT - 1
                                                    GO TO 60

```

```

C
C   ***MODE=1,2***   ADD JADD BASIS FCTNS, I.E., FOR IO=1,JADD,
C                   CONSTRUCT FCT ILAST WITH ONE MORE KNOT, VIZ.
C                   XKNOT=ADDXI(IO), THAN THE PREVIOUS LAST FCT,

```

```

C          ORTHONORMALIZE IT OVER ALL PREVIOUS FCTNS, THEN
C          COMPUTE THE COORDINATE BC(ILAST) OF U WRTO IT,
C          SUBTRACT OUT ITS COMPONENT FROM UERROR.
50 IF (JADD.LE.0)          GO TO 61
51 DO 52 IO=1,JADD
    XKNOT = ADDXI(IO)
    CALL NUBAS
    DO 52 L=1,LX
        UERROR(L) = UERROR(L) - BC(ILAST)*FCT(L,ILAST)
52 continue
C
60 FXDKNT= DOT(31,2)/XSCALE
    ERBUT1 = FXDKNT
    KNOTSV = KNOT
61 MODE3 = .FALSE.
    IF (IPRINT.EQ.0)          RETURN
C          VARIOUS PRINTING IS DONE DEP ON THE IPRINT = IFIX(CHANGE)
C          GO TO (70,80,90),IPRINT
C
C          COMPUTE COEFFICIENTS OF B.A. AND PRINT
C          ****          BEST APPROXIMATION PRINTOUT          ****
C          FORMAT IS
C          KNOTS XI(J)          CUBIC COEFFICIENTS P(I,J) IN
C          INTERVAL (XI(J), XI(J+1))
C          ERROR CURVE (SCALED)
C
C          THE FOLLOWING FORTRAN CODE FINDS VALUES AT X OF THE
C          APPROXIMATION FROM THIS OUTPUT----
C          I=LXI
C          1 A=X-XI(I)
C          IF(A .ge. 0.) goto 4
C          I=I-1
C          IF(I .gt. 0) goto 1
C          I=1
C          4 V=P(1,I)+A*(P(2,I)+A*(P(3,I)+A*(P(4,I))))
C
70 WRITE(6,610)

```



```

DO 72 I=1,KNOT
  ILOC = IORDER(I)
  DO 72 L=1,2
    SUM = 0.DO
    DO 71 J=1,ILAST
      SUM = SUM + BC(J)*VORD(J,ILOC,L)
71    continue
      VORDL(I,L) = SUM
72 continue
  CALL EVAL
  DO 73 I=1,INTERV
    WRITE(6,620) I,XIL(I)
    WRITE (6,630) (J,COEFL(I ,J),J=1,4)
73 continue
  WRITE (6,620) KNOT,XIL(KNOT)
610 FORMAT(12X,'KNOTS',22X,'CUBIC COEFFICIENTS'//)
620 FORMAT(5X,'XI(',I2,') =', F12.6)
630 FORMAT(37X,'C(',I1,') =',E16.6)
C
C    **COMPUTE L2, L1, MAX ERRORS AND PRINT
80 ERRL2 = SQRT(FXDKNT)
  ERRL1 = 0.
  ERRL99= 0.
  DO 82 L=1,LX
    DIF = ABS(UERROR(L)*WEIGHT(L))
    IF(ERRL99.GT.DIF) GO TO 81
    LMAX = L
    ERRL99 = DIF
81  ERRL1 = ERRL1+ DIF
82 CONTINUE
  ERRL1 = ERRL1/FLOAT(LX)
  WRITE(6,623) ERRL2, ERRL1, ERRL99,XX(LMAX)
C    *** THE FOLLOWING CARD IS TEMPORARY
  GO TO (90,96,96),IPRINT
C
C    ** SCALE ERROR CURVE AND PRINT
90 IE = 0

```

```

SCALE = 1.
IF (ERRL99.GE.10.)          GO TO 92
DO 91 IE=1,9
    SCALE = SCALE*10.
    IF (ERRL99*SCALE.GE.10.)  GO TO 92
91 CONTINUE
92 DO 93 L=1,LX
    SERROR(L) = UERROR(L)*SCALE
93 continue
                                GO TO (94,95,95),IPRINT
94 WRITE (6,621) IE,(L,XX(L),FCTL(L),SERROR(L),L=1,LX)
                                GO TO 96
95 WRITE (6,622) IE,(L,XX(L),SERROR(L),L=1,LX)
96                                RETURN
621 FORMAT(' '//15X,'APPROXIMATION AND SCALED ERROR CURVE'/8X,
*'DATA POINT',7X,'APPROXIMATION',3X,'DEVIATION X 10E+',I1/
*(1X,I4,F16.8,F16.8,F17.6))
622 FORMAT(' //' ' ERROR CURVE'/ 8X, 'DATA POINT', 23X,
1'DEVIATION X 10E+',I1/( 1X,I4,F16.8,16X,F17.6))
623 FORMAT(' '///10X,'LEAST SQUARE ERROR =',E20.6/
1          10X,'AVERAGE ERROR      =',E20.6/
2          10X,'MAXIMUM ERROR       =',E20.6,' AT',F12.6///)
END
C
C*****
C
SUBROUTINE INTERP
C
C      COMPUTE THE SLOPES VORDL(I,2), I=2,KNOT-1 AT INTERIOR
C      KNOTS OF CUBIC SPLINE FOR GIVEN VALUES VORDL(I,1),I=1,KNOT,
C      AT ALL THE KNOTS AND GIVEN BOUNDARY DERIVATIVES
IMPLICIT NONE
real COEFL(27,4),D(28),DIAG(28),FCTL(100),G,UERROR(100),
1  VORDL(28,2),XIL(28)
integer INTERV,KNOT,LMAX,M,NJ
COMMON/ OUTPUT /UERROR,FCTL,XIL,COEFL,VORDL,KNOT,LMAX,INTERV
DATA DIAG(1),D(1)/1.,0./

```

```

DO 10 M=2,KNOT
  D(M) = XIL(M) - XIL(M-1)
  DIAG(M) = (VORDL(M,1)-VORDL(M-1,1))/D(M)
10 continue
DO 20 M=2,INTERV
  VORDL(M,2) = 3.*(D(M)*DIAG(M+1) + D(M+1)*DIAG(M))
  DIAG(M) = 2.*(D(M)+D(M+1))
20 continue
DO 30 M=2,INTERV
  G = -D(M+1)/DIAG(M-1)
  DIAG(M) = DIAG(M) + G*D(M-1)
  VORDL(M,2) = VORDL(M,2) + G*VORDL(M-1,2)
30 continue
  NJ = KNOT
DO 40 M=2,INTERV
  NJ = NJ - 1
  VORDL(NJ,2) = (VORDL(NJ,2) - D(NJ)*VORDL(NJ+1,2))/DIAG(NJ)
40 continue

                                RETURN

END

```

```

C
C*****
C
  FUNCTION DOT (M,INDEX)
C      COMPUTE INNER PRODUCT OF FCT M WITH FCT ILAST (INDEX=1) OR
C      UERROR (INDEX=2)
  IMPLICIT NONE
  real ADDXI(26),BC(30),COEFL(27,4),DOT,
1    FCT(100,30),FCTL(100),
2    TREND(100),
3    U(100),UERROR(100),VORD(30,28,2),VORDL(28,2),
4    XIL(28),XX(100)
  integer ILAST,INDEX,INTERV,
1    JADD,KNOT,L,LMAX,LX,M,MODE
  DOUBLE PRECISION TRPZWT(100),G(100),SUM
  COMMON / WANDT / TREND,TRPZWT
  COMMON/ INPUT /LX,XX,U,JADD,ADDXI,MODE

```

```

COMMON/ OUTPUT /UERROR,FCTL,XIL,COEFL,VORDL,KNOT,LMAX,INTERV
COMMON/ BASIS /FCT,VORD,BC,ILAST
                                GO TO (10,30),INDEX
10 IF (M.EQ.ILAST)              GO TO 20
    DO 11 L=1,LX
        G(L) = FCT(L,M)*FCTL(L)
11 continue
                                GO TO 80
20 DO 21 L=1,LX
        G(L) = FCTL(L)*FCTL(L)
21 continue
                                GO TO 80
30 IF (M.EQ.31)                 GO TO 40
    DO 31 L=1,LX
        G(L) = FCTL(L)*UERROR(L)
31 continue
                                GO TO 80
40 DO 41 L=1,LX
        G(L) = UERROR(L)*UERROR(L)
41 continue
C
C EFFICIENTLY PROGRAMMED DOUBE PRECISION ACCUMULATION OF SCALAR
C PRODUCTS IS CALLED FOR HERE.  AT PURDUE, WE USE
C     D = ARITH1(C,N,A,IA,B,IB)
C WHICH RETURNS THE VALUE OF
C     D = C - SUM(A(1+J*IA) * B(1+J*IB), J=0,...,N-1)
C
80 SUM = 0.DO
    DO 81, L=1,LX
        SUM = SUM + G(L)*TRPZWT(L)
81 continue
    DOT = SUM
                                RETURN
END
C
C*****
C

```

```

SUBROUTINE EVAL
C      COMPUTE POL. COEFF COEFL(I,K) OF FCT ILAST FROM VORDL,
C      THEN COMPUTE FCTL(L) = (FCT ILAST)*TREND AT XX(L),L=1,LX
C
      IMPLICIT NONE
      real ADDXI(26),COEFL(27,4),DUM1,DUM2,DX,
1      FCTL(100),
2      TREND(100),TRPZWT(100),
3      U(100),UERROR(100),VORDL(28,2),
4      XIL(28),XX(100)
      integer I,INTERV,ISWTCH,
1      J,JADD,KNOT,L,LMAX,LX,MODE
      COMMON / WANDT / TREND,TRPZWT
      COMMON/ INPUT /LX,XX,U,JADD,ADDXI,MODE
      COMMON/ OUTPUT /UERROR,FCTL,XIL,COEFL,VORDL,KNOT,LMAX,INTERV
      DO 10 I=1,INTERV
          COEFL(I,1) = VORDL(I,1)
          COEFL(I,2) = VORDL(I,2)
          DX = XIL(I+1) - XIL(I)
          DUM1 = (VORDL(I+1,1)-VORDL(I,1))/DX
          DUM2 = VORDL(I,2)+VORDL(I+1,2)-2.*DUM1
          COEFL(I,3) = (DUM1-DUM2-VORDL(I,2))/DX
          COEFL(I,4) = DUM2/DX/DX
10 continue
C
      J = 1
      ISWTCH = 1
      DO 20 L=1,LX
          GO TO (11,13),ISWTCH
11      IF (J.EQ.INTERV)          GO TO 12
          IF (XX(L).LT.XIL(J+1))  GO TO 13
          J = J + 1
          GO TO 11
12      ISWTCH = 2
13      DX = XX(L) - XIL(J)
          FCTL(L) = (COEFL(J,1)+DX*(COEFL(J,2)+DX*(COEFL(J,3)
*                               +DX*COEFL(J,4))))*TREND(L)

```

20 continue

RETURN

END

C

C\*\*\*\*\*

C

SUBROUTINE NUBAS

IMPLICIT NONE

real ADDXI(26),BC(30),C,COEF(381,4),COEFL(27,4),

1 DOT,DX,FCT(100,30),FCTL(100),

2 ONEOVC,TEMP(30),

3 U(100),UERROR(100),VORD(30,28,2),VORDL(28,2),

4 XI(381),XIL(28),XKNOT,XX(100)

integer I,IBOUND,ICLAST,ID,ILAST,ILM1,ILOC,INSERT,INSIRT(30),

1 INTERV,IO,IORDER(28),JADD,K,KNOT,L,LMAX,LX,MODE

COMMON/ INPUT /LX,XX,U,JADD,ADDXI,MODE

COMMON/ OUTPUT /UERROR,FCTL,XIL,COEFL,VORDL,KNOT,LMAX,INTERV

COMMON/ BASIS /FCT,VORD,BC,ILAST

COMMON/ LASTB /IORDER,INSIRT,XKNOT

C COEF(IC,.) CONTAINS THE POL COEFFICIENTS OF FCT M FOR INTER-

C VAL TO THE RIGHT OF XI(IC), IC=ICM,ICM+M-3,

C WITH ICM = M\*(M-7)/2 + 10 (WITH OBVIOUS MODS FOR M.LE.4)

C THE FCT ILAST (TO BE) INTRODUCED LAST, HAS ITS VALUES AT THE

C THE POINTS XX(L) IN FCTL(L), HAS FIRST INDEX ICLAS

C IN COEF AND XI, HAS ADDITIONAL KNOT XKNOT, THE KNOT KNOTS

C FOR IT ARE CONTAINED, IN INCREASING ORDER, IN XIL,ITS COR-

C RESPONDING ORDS AND SLOPES ARE IN VORDL, THE KNOT JUST INTRO

C DUCED HAS INDEX INSERT IN XIL,INSERT IS SAVED IN INSIRT(ILAS

C FOR POSSIBLE REPLACEMENT OF KNOTS LATER ON (SEE MODE=2,3).

logical REPEAT

REPEAT = .FALSE.

IF (MODE.GT.0) GO TO 8

C-----\*\*CONSTRUCT FCT ILAST FOR ILAST.LE.4

XI(ILAST) = XIL(1)

ICLAST = ILAST

ILM1 = ILAST-1

IF (ILAST.GT.2) GO TO 7

```

        IF (ILAST.EQ.2)                GO TO 6
C      FIRST BASIS FCT IS A CONSTANT
        VORDL(1,1) = 1.
        VORDL(2,1) = 1.
        VORDL(1,2) = 0.
        VORDL(2,2) = 0.
                                           GO TO 67
C      SECOND BASIS FCT IS A STRAIGHT LINE
6     VORDL(2,2) = VORDL(1,1)/(XIL(2) - XIL(1))*2.
        VORDL(1,2) = -VORDL(2,2)
C
7     VORDL(2,1) = - VORDL(2,1)
        VORDL(2,2) = - VORDL(2,2)
                                           GO TO 59
C-----
        8                                GO TO (10,10,140),MODE
C-----***SET UP CONSTANTS DEP.ON ILAST. INSERT NEW KNOT INTO XIL
C      AND UPDATE VORD FOR FCT M,M=1,ILAST-1
10    KNOT = KNOT + 1
        ILAST = KNOT + 2
        ICLAST = ILAST*(ILAST-7)/2 + 10
        ILM1 = ILAST-1
        INTERV = KNOT - 1
        DO 11 INSERT=2,INTERV
        IF (XKNOT.LT.XIL(INSERT))      GO TO 12
11    CONTINUE
                                           GO TO 95
12    IF (XKNOT.LE.XIL(INSERT-1))     GO TO 95
        IO = KNOT
        DO 13 L=INSERT,INTERV
        IO = IO - 1
        XIL(IO+1) = XIL(IO)
13    IORDER(IO+1) = IORDER(IO)
        IORDER(INSERT) = KNOT
                                           go to 14
140   insert = insirt(ilast)
14    XIL(INSERT) = XKNOT

```

```

DX = XKNOT - XIL(1)
DO 15 I=1,4
VORD(I,KNOT,1)=COEF(I,1)+DX*(COEF(I,2)+DX*(COEF(I,3)
*
+DX*COEF(I,4)))
15 VORD(I,KNOT,2)=COEF(I,2)+DX*(2.*COEF(I,3)+DX*3.*COEF(I,4))
IF(ILM1.LT.5) GO TO 20
ID = 4
IBOUND = 4
DO 19 I=5,ILM1
ID = ID + I - 4
IBOUND = IBOUND + I - 3
17 IF (ID.EQ.IBOUND) GO TO 18
IF (XKNOT.LT.XI(ID+1)) GO TO 18
ID = ID + 1
GO TO 17
18 DX = XKNOT - XI(ID)
VORD(I,KNOT,1)=COEF(ID,1)+DX*(COEF(ID,2)+DX*(COEF(ID,3)
*
+DX*COEF(ID,4)))
19 VORD(I,KNOT,2)=COEF(ID,2)+DX*(COEF(ID,3)*2.+DX*3.*COEF(ID,4))
C-----
C-----DEFINE LAST BASIS FUNCTION
20 CONTINUE
GO TO (30,40,50),MODE
C *** MODE=1 *** ADD ILAST-TH BASIS FUNCTION. CONSTRUCT FROM FCT
C ILAST-1 BY REFLECTING THE PART OF THE LATTER TO
C THE RIGHT OF XKNOT ACROSS THE X-AXIS, THEN INTER
C POLATING. THIS SHOULD INDUCE ONE MORE OSCILLATIO
C N IN FCT ILAST THAN IN FCT ILAST-1
C
29 MODE = 1
30 VORDL(1,2) = 0.
DO 31 K=1,KNOT
31 VORDL(K,1) = MAX(0.,XIL(K)-XKNOT)**3
VORDL(KNOT,2) = 3.*(XIL(KNOT)-XKNOT)**2
GO TO 55
C
C *** MODE=2 *** REPLACE FCT ILAST BY INTERPOLATING IT AT THE

```



```

C          CURRENT SET OF KNOTS. IF FCT ILAST HAS NOT BEEN
C          PREVIOUSLY DEF (INSIRT(ILAST)=0)(SEE 9 ABOVE,
C          ALSO MAIN AT 10)) SET MODE=1,PROCEED IN THAT MOD
C
40 IF (INSIRT(ILAST).EQ.0)          GO TO 29
   VORDL(1,1)=VORD(ILAST,1,1)
   VORDL(1,2)=VORD(ILAST,1,2)
   ID = ICLAST
   IBOUND = ICLAST + ILAST - 4
   DO 43 K=2,INTERV
41 IF (ID.EQ.IBOUND)          GO TO 42
   IF (XIL(K).LT.XI(ID+1))    GO TO 42
   ID = ID +1
                               GO TO 41
42 DX = XIL(K)- XI(ID)
43 VORDL(K,1) = COEF(ID,1)+DX*(COEF(ID,2)+DX*(COEF(ID,3)
*                               +DX*COEF(ID,4)))
   VORDL(KNOT,1)=VORD(ILAST,2,1)
   VORDL(KNOT,2)=VORD(ILAST,2,2)
                               GO TO 55
C
C          *** MODE=3 *** CHANGE FCT ILAST BY CHANGING JUST THE KNOT INTRO
C          DUCED LAST
C
50 ID = ICLAST + INSERT - 1
   DX = XKNOT - XI(ID)
   XI(ID) = XKNOT
   IF (DX.GE.0.)          GO TO 51
   ID = ID - 1
   DX = XKNOT - XI(ID)
51 VORDL(INSERT,1) = COEF(ID,1) +DX*(COEF(ID,2)+DX*(COEF(ID,3)
*                               +DX*COEF(ID,4)))
C
C          *** INTERPOLATE
55 CALL INTERP
                               GO TO (57,57,59),MODE
57 ID = ICLAST - 1

```

```

DO 56 IO=1,INTERV
  ID = ID + 1
56 XI(ID) = XIL(IO)
  INSIRT(ILAST) = INSERT
C-----
C-----*** ORTHONORMALIZE FCT ILAST OVER PREVIOUS (ORTHONORMAL) SET
C          THEN COMPUTE THE COMPONENT BC(ILAST) OF UERROR WRTO IT
C          FINALLY,STORE THE VARIOUS REPRESENTATIONS OF FCT ILAST
C
59 CALL EVAL
  TEMP(ILAST) = SQRT(DOT(ILAST,1))
  IF (REPEAT .AND. ABS(1.-TEMP(ILAST))) .GT. .5) GO TO 65
  DO 60 I=1,ILM1
  TEMP(I) = DOT(I,1)
  DO 69 L=1,LX
69   FCTL(L) = FCTL(L) - TEMP(I)*FCT(L,I)
  DO 61 K=1,KNOT
  ILOC = IORDER(K)
  DO 61 L=1,2
61   VORDL(K,L) = VORDL(K,L) - TEMP(I)*VORD(I,ILOC,L)
60 CONTINUE
67 CALL EVAL
  C = SQRT(DOT(ILAST,1))
  WRITE (6,667)ILAST,(TEMP(I),I=1,ILAST),C
667 FORMAT(I3,7E11.3/(7X,7E11.3))
  IF(ILAST .GT. 1 .AND. REPEAT .AND. ABS(1.-C) .GT. .5) GO TO 65
  IF (C+TEMP(ILAST) .LE. TEMP(ILAST)) GO TO 65
  ONEOVC = 1./C
                                     GO TO 68
65 ONEOVC = 0.
68 BC(ILAST) = DOT(ILAST,2)*ONEOVC
  DO 62 K=1,KNOT
  ILOC = IORDER(K)
  DO 62 L=1,2
  VORDL(K,L) = VORDL(K,L)*ONEOVC
62 VORD(ILAST,ILOC,L) = VORDL(K,L)
  IF (ONEOVC .EQ. 0. .OR. ILAST .EQ. 1

```

```

*          .OR. C .GE. 1.E-2*TEMP(ILAST)) GO TO 152
REPEAT = .TRUE.
CALL INTERP
                                GO TO 59
152 CONTINUE
    ID = ICLAST - 1
    DO 63 IO=1,INTERV
    ID = ID + 1
    DO 63 L=1,4
63 COEF(ID,L) = COEFL(IO,L)*ONEOVC
    DO 64 L=1,LX
64 FCT(L,ILAST) = FCTL(L)*ONEOVC
C-----
                                RETURN
C
C   *** THIS OUTPUT INDICATES A FAILURE CONDITION ***
95 WRITE (6,950) XKNOT,ILAST
950 FORMAT (15H   *** NEW KNOT,E20.8,13H FOR FUNCTION,I3,50H OUT OF BO
    *UNDS OR COINCIDENT WITH A PREVIOUS KNOT./36H   *** EXECUTION CANNO
    *T BE CONTINUED)
                                STOP
C
    END
C
C*****TREND AND WEIGHT FUNCTIONS*****
C
    FUNCTION T(Z)
    real T,Z
    T = 1.
    RETURN
    END
C
    FUNCTION W(Z)
    W = 1.
    RETURN
C
    END

```